

# CO444: Algebraic Graph Theory

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# Chapter 1

## Introduction

### 1.1 Graphs

Recall that a (simple) *graph* is a pair  $X = (V, E)$  consisting of a vertex set and an edge set.

**Definition 1.1.1 (Isomorphism)**

An isomorphism between graphs  $X, Y$  is a function  $f : V(X) \rightarrow V(Y)$  such that  $uv \in E(X) \iff f(u)f(v) \in E(Y)$ .

If an isomorphism exists between  $X$  and  $Y$ , we say  $X$  and  $Y$  are *isomorphic* and write

$$X \cong Y.$$

Some important classes of graphs are complete graphs, bipartite graphs, empty graphs ( $E = \emptyset$ ), the null graph ( $V = \emptyset$ ), multigraphs, simple graphs, directed graphs, finite graphs, and infinite graphs.

We will focus mainly on finite simple graphs, with some exploration of finite multigraphs.

### 1.2 Subgraphs

Recall that a *subgraph*  $Y$  of the graph  $X$  is a graph such that  $V(Y) \subseteq V(X)$ ,  $E(Y) \subseteq E(X)$ . Some important classes of subgraphs are spanning subgraphs and induced subgraphs. Some other examples of graphs are cliques, independent sets, paths, cycles, and spanning trees.

## 1.3 Automorphisms

### Definition 1.3.1 (Automorphism)

An automorphism of the graph  $X$  is an isomorphism  $f : X \rightarrow X$ .

We write  $\text{Aut}(X)$  to denote the set of all automorphisms of  $X$ .

If  $\text{Sym}(V(X))$  is the permutation group of the vertices of  $X$ , then it is clear that  $\text{Aut}(X) \subseteq \text{Sym}(V(X))$ . We will sometimes write  $\text{Sym}(n)$  to denote  $\text{Sym}(V(X))$  as we can label the vertices using  $[n]$ .

### Proposition 1.3.1

$\text{Aut}(X)$  is a subgroup of  $\text{Sym}(V(X))$ .

#### Proof

It is a group and a subset.

For  $h \in \text{Sym}(V(X))$  and  $v \in V(X)$ , we let  $v^g := g(v)$  denote the image of  $v$  under  $g$ . Similarly, for  $S \subseteq V(X)$ , we let  $S^g := \{v^g : v \in S\}$ .

Suppose  $Y$  is a subgraph of  $X$ . We let  $Y^g$  be the graph where

$$\begin{aligned} V(Y^g) &:= V(Y)^g \\ E(Y^g) &:= \{u^g v^g : uv \in E(Y)\} \end{aligned}$$

It is not hard to see that  $Y^g$  is another subgraph of  $X$  that is in fact isomorphic to  $Y$ .

### Lemma 1.3.2

For  $g \in \text{Aut}(X)$  and  $v \in V(X)$ ,

$$\deg_X(v) = \deg_{X^g}(v^g).$$

Hence automorphisms permute the vertices of equivalent degrees.

#### Proof

Let  $Y$  be the subgraph induced by  $N(v) + v$ . The result follows as  $Y^g \cong Y$ .

Let  $d(u, v)$  denote the length of the shortest path between  $u, v$ .



**Lemma 1.3.3**

For  $g \in \text{Aut}(X)$  and  $u, v \in V(X)$ ,

$$d_X(u, v) = d_{X^g}(u^g, v^g).$$

Hence automorphisms preserve distances.

**Proof**

Using  $g, g^{-1}$ , it is not hard to see that there is a  $uv$ -path of length  $\ell$  in  $X$  if and only if there is a  $u^g v^g$ -path of length  $\ell$  in  $X^g$ .

## 1.4 Homomorphisms

**Definition 1.4.1 (Homomorphism)**

A homomorphism between graphs  $X, Y$  is a function  $f : V(X) \rightarrow V(Y)$  such that

$$xy \in E(X) \implies f(x)f(y) \in E(Y).$$

If such a homomorphism exists, we say  $X$  is *homomorphic* to  $Y$ .

This is a much more relaxed condition than an isomorphism.

**Example 1.4.1**

Given any bipartite graph  $X$  with bipartition  $A \cup B$ ,  $X$  is homomorphic to the graph consisting of a single edge. Indeed, simply map  $A$  to one endpoint and  $B$  to another.

**Example 1.4.2**

Any subgraph is homomorphic to its supergraph through the identity map.

Recall the definition of a valid *graph colouring*. Let  $\chi(X)$  denote the *chromatic number* of  $X$ .

**Lemma 1.4.3**

$\chi(X)$  is equal to the minimum integer  $r$  such that there exists a homomorphism from  $X$  to  $K_r$ .

**Proof**

First we claim that  $\chi(X) \geq r$ . Indeed, let  $c : V(X) \rightarrow [\chi(X)]$  be a  $\chi(X)$ -colouring of  $X$ . If we take  $[\chi(X)]$  to be the vertices of  $K_{\chi(X)}$ , then  $c$  is a homomorphism from  $X \rightarrow K_{\chi(X)}$ ! Hence  $\chi(X) \geq r$ .

Conversely, suppose  $f : X \rightarrow K_r$  is a homomorphism. Then  $f^{-1}(\{i\})$  induces an independent set for all  $i \in [r]$ . Thus  $X$  necessarily has an  $r$ -colouring. It follows that  $\chi(X) \leq r$ .

Remark that we have in fact proven that the set of homomorphisms  $X \rightarrow K_r$  is in fact the set of  $r$ -colorings of  $X$  with colours  $[r]$ !

**Definition 1.4.2 (Retraction)**

A retraction is a homomorphism from  $X$  to a subgraph  $Y$  of  $X$  such that the restriction  $f|_Y$  is the identity map.

If such a retraction exists, we say that  $Y$  is a *retract* of  $X$ .

**Example 1.4.4**

Any clique  $Y$  of  $X$  that is the size of  $\chi(X)$  is a retract of  $X$ .

## 1.5 Examples of Graphs

### 1.5.1 Circulant Graphs

Consider the cycle  $C_n$ . Its automorphism group is exactly the dihedral group on  $n$  elements.

We can view circulant graphs as a generalization of cycles.

**Definition 1.5.1 (Circulant Graph)**

Let the vertex set be  $\mathbb{Z}_n$ . Choose a subset  $C \subseteq \mathbb{Z}_n - 0$  that is closed under inverses. Define

$$E(X) := \{ij : i - j \in C\}.$$

We write  $X = X(\mathbb{Z}_n, C)$  to denote the circulant graph parametrized by  $\mathbb{Z}_n, C$ .

Let  $g \in \text{Aut}(X)$  be the cyclic shift and  $h \in \text{Aut}(X)$  be the map sending  $i \mapsto -i$ . Furthermore, put  $R$  as the cyclic subgraph generated by  $g$ . It follows that  $R, hR$  are different cosets of

$\text{Aut}(X)$  and

$$|\text{Aut}(X)| \geq 2n.$$

## 1.5.2 Johnson Graphs

### Definition 1.5.2 (Johnson Graph)

Let  $v \geq k \geq i$  be integers. The Johnson graph  $J = J(n, k, i)$  is the graph with vertex set

$$V(J) := \{S \subseteq [v] : |S| = k\}$$

and edge set

$$E(J) := \{ST : |S \cap T| = i\}.$$

### Proposition 1.5.1

$J(v, k, i)$  is  $d$ -regular for

$$d = \binom{k}{i} \binom{v-k}{k-i}.$$

### Proof

regularity is clear. It remains only to compute the common degree of all vertices.

Fix a vertex  $S$ . Fix a subset of  $S' \subseteq S$  with size  $k$ . There are  $\binom{v-k}{k-i}$  vertices  $T$  such that

$$S \cap T = S'.$$

The result follows by letting  $S'$  range over all  $S$ -subsets of size  $k$ .

### Lemma 1.5.2

$$J(v, k, i) \cong J(v, v-k, v-2k+i).$$

### Proof

The map  $f(S) := \bar{S}$  yields an isomorphism.

The special case of  $J(v, k, 0)$  are the *Kneser Graphs* and  $J(5, 2, 0)$  is the *Peterson Graph*.

**Lemma 1.5.3**

$\text{Aut}(J(v, k, i))$  contains a subgroup isomorphic to  $\text{Sym}(v)$ .

**Proof**

Let  $g \in \text{Sym}(v)$ . Define the map  $\sigma_g : V(J(v, k, i)) \rightarrow V(J(v, k, i))$  by

$$S \mapsto S^g.$$

It is easy to see that  $|S \cap T| = |S^g \cap T^g|$ , so  $\sigma_g \in \text{Aut}(J(v, k, i))$ . But then

$$\{\sigma_g : g \in \text{Sym}(v)\} \subseteq \text{Aut}(J(v, k, i))$$

and  $\{\sigma_g : g \in \text{Sym}(v)\} \cong \text{Sym}(v)$  as required.

Remark that we typically have  $\text{Aut}(J(v, k, i)) \cong \text{Sym}(v)$  but it is NOT always true.

**1.5.3 Line Graphs****Definition 1.5.3 (Line Graph)**

The line graph of  $X$  is the graph  $L(X)$  with vertex set

$$V(L(X)) := E(X)$$

and edge set

$$E(L(X)) := \{ef : e, f \text{ share an endpoint}\}.$$

In general, if  $X \cong Y$ , then  $L(X) \cong L(Y)$ . But the converse need not hold! Indeed,  $L(C_3) = L(K_{1,3})$  but  $C_3 \neq K_{1,3}$ .

The converse is true however, if the minimum degrees of  $X, Y$  are at least 4.

It is also interesting to observe that not all graphs are line graphs. For example,  $K_{1,3}$  is not a line graph since no matter how we distribute the outside vertices according to the endpoints of the central vertex, there are not enough edges for  $K_{1,3}$  to be a line graph.

# Chapter 2

## Group Actions

### 2.1 Group Actions on Graphs

Recall the definition of a group.

**Definition 2.1.1 (Homomorphism)**

Let  $G, H$  be groups. A map  $f : G \rightarrow H$  is a homomorphism if

$$f(xy) = f(x)f(y)$$

for all  $x, y \in G$ .

The *kernel* of a homomorphism is

$$\ker f := f^{-1}(1).$$

**Definition 2.1.2 (Group Action)**

Let  $G$  be a group and  $V$  a set. A homomorphism  $f : G \rightarrow \text{Sym}(V)$  is an action of  $G$  on  $V$ .

We say  $G$  *acts* on  $V$ .

If in addition,  $f^{-1}(1) = 1$ , then we say  $f$  is *faithful*.

Suppose  $G$  acts on  $V$ . Let  $g \in G$ . We write  $x^g$  to denote the image of  $x$  under  $f(g)$ . Similarly, for  $S \subseteq V$ , we write

$$S^g := \{x^g : x \in S\}.$$

**Definition 2.1.3 (Invariant)**

Let  $G$  act on  $V$ .  $S \subseteq V$  is  $G$ -invariant if

$$S^g = S$$

for all  $g \in G$ .

**Definition 2.1.4 (Orbit)**

Let  $G$  act on  $V$ . The orbit of  $x \in V$  is

$$x^G := \{x^g : g \in G\}.$$

It is well-known that  $B$  is partitioned into disjoint orbits and each such orbit is  $G$ -invariant.

## 2.1.1 Orbits & Stabilizers

**Definition 2.1.5**

Let  $G$  be a permutation group acting on  $V$ . Fix  $x \in V$ .

The stabilizer of  $x$  is

$$G_x := \{g \in G : x^g = x\}.$$

It is not hard to see that  $G_x$  is a subgroup of  $G$ .

**Lemma 2.1.1**

Let  $G$  be a permutation group acting on  $V$  and  $S$  an orbit of  $G$ .

Suppose  $x, y \in S$ . Then  $H := \{h \in G : x^h = y\}$  is a left coset of  $G_x$ .

Conversely, if  $H$  is a left coset of  $G_x$ , then  $x^h = x^{h'}$  for all  $h, h' \in H$ .

**Proof**

Since  $G$  is transitive on  $S$ , there is some  $g \in G$  such that

$$x^g = y.$$

But then for any  $h \in H$  such that  $x^h = y = x^g$ ,

$$\begin{aligned}x^h &= x^g \\x^{g^{-1}h} &= x \\g^{-1}h &\in G_x \\h &\in gG_x.\end{aligned}$$

Hence  $H$  is indeed a left coset of  $G_x$ .

Now, consider the left coset  $gG_x$  for some  $g \in G$ . Pick  $g\sigma, g\sigma' \in gG_x$  for some  $\sigma, \sigma' \in G_x$ . We have

$$\begin{aligned}x^{g\sigma} &= x^{g\sigma'} \\&= x^g \\&= x^{g\sigma'} \\&= x^{h'}\end{aligned}$$

as desired.

### Lemma 2.1.2 (Orbit-Stabilizer)

Let  $G$  be a permutation group acting on  $V$  and  $x \in V$ . Then

$$|G_x| \cdot |x^G| = |G|.$$

#### Proof

Let  $\mathcal{H}$  be the collection of left cosets of  $G_x$  and  $f : x^G \rightarrow \mathcal{H}$  be the map

$$f(y) := \{g \in G : x^g = y\}$$

for all  $y \in x^G$ . By the previous lemma,  $f$  is bijective.

It follows that  $|x^G| = |\mathcal{H}|$ . Since the left cosets of  $G_x$  partition  $G$ , we must have

$$|G| = |G_x| \cdot |\mathcal{H}| = |G_x| \cdot |x^G|.$$

## 2.2 Burnside's Lemma

What is the relationship between  $G_x, G_y$  for  $y \in x^G$ ?

**Definition 2.2.1 (Conjugate)**

Let  $g, h \in G$ .  $g$  is conjugate to  $h$  if there is some  $\sigma \in G$  where

$$g = \sigma^{-1}h\sigma.$$

Conjugacy is an equivalence relation so  $G$  is partitioned into conjugacy classes.

For  $g \in G$ , let  $\tau_g : G \rightarrow G$  be the map sending

$$h \mapsto ghg^{-1}$$

for all permutations  $h \in G$ . The conjugacy classes of  $G$  are the orbits of  $G$  under  $\{\tau_g : g \in G\}$ .

Let  $H$  be a subgroup of  $G$ . Then  $gHg^{-1}$  is a subgroup of  $G$ . Moreover, it is possible to show that  $gHg^{-1} \cong H$  and we say  $gHg^{-1}$  is conjugate to  $H$ .

**Lemma 2.2.1**

Let  $G$  be a permutation group acting on  $V$  and  $x \in V$ . Then for all  $g \in G$ ,

$$gG_xg^{-1} = G_{x^g}.$$

Thus if  $y \in x^G$ , then  $G_x, G_y$  are conjugate.

**Proof**

$gG_xg^{-1} \subseteq G_{x^g}$  Let  $y = x^g$  so that  $y^{g^{-1}} = x$ . For any  $ghg^{-1} \in gG_xg^{-1}$ , we have

$$\begin{aligned} y^{ghg^{-1}} &= x^{gh} \\ &= x^g \\ &= y. \end{aligned}$$

Hence  $gG_xg^{-1} \subseteq G_{x^g}$ .

$G_{x^g} \subseteq gG_xg^{-1}$  For any  $h \in G_y$ , we have

$$\begin{aligned} x^{g^{-1}hg} &= y^{g^{-1}h} \\ &= y^{g^{-1}} \\ &= x. \end{aligned}$$

It follows that  $g^{-1}hg \in G_x$  and  $h \in gG_xg^{-1}$ .



Let  $G$  be a permutation group acting on  $V$  and  $g \in G$ . We write

$$\text{fix}(g) := \{v \in V : v^g = v\}.$$

**Lemma 2.2.2 (“Burnside”)**

Let  $G$  be a permutation group acting on  $V$ , then the number of orbits of  $G$  is given by

$$\frac{1}{|G|} \sum_{g \in G} |\text{fix}(g)|.$$

**Proof**

Let  $\Lambda := \{(g, x) : g \in G, x \in \text{fix}(g)\}$ . We will apply a double counting argument.

$$\begin{aligned} \sum_{g \in G} |\text{fix}(g)| &= |\Lambda| \\ &= \sum_{x \in V} |G_x| \\ &= \sum_{x \in V} \frac{|G|}{|x^G|} && \text{Orbit-Stabilizer Lemma} \\ &= |G| \cdot (\text{number of orbits}). \end{aligned}$$

## 2.3 Asymmetric Graphs

**Definition 2.3.1 (Asymmetric)**

A graph  $X$  is asymmetric if its automorphism group  $\text{Aut}(X) = \{1\}$  is trivial.

Let  $\mathcal{G}_n$  denote the set of graphs with vertex set  $[n]$ . Then we write

$$\text{Iso}(X) := \{Y \in \mathcal{G}_n : Y \cong X\}$$

to denote the isomorphism class of  $X$ .

**Lemma 2.3.1**

For all  $X \in \mathcal{G}_n$ ,

$$|\text{Iso}(X)| = \frac{n!}{|\text{Aut}(X)|}.$$

**Proof**

Let  $G = \text{Sym}(n)$ . For  $g \in G$ , let  $\tau_g : \mathcal{G}_n \rightarrow \mathcal{G}_n$  be the map sending

$$X \mapsto X^g.$$

The group  $H := \{\tau_g : g \in G\}$  acts on  $\mathcal{G}_n$  and is isomorphic to  $G$ . By the Orbit-Stabilizer Lemma

$$\begin{aligned} n! &= |G| \\ &= |H| \\ &= |X^H| \cdot |H_X| \\ &= |\text{Iso}(X)| \cdot |\text{Aut}(X)|. \end{aligned}$$

**Lemma 2.3.2**

The set  $\mathcal{H}$  of all isomorphism classes of  $\mathcal{G}_n$  satisfies

$$|\mathcal{H}| = (1 + o(1)) \frac{2^{\binom{n}{2}}}{n!}.$$

**Proof**

Consider the group

$$\mathcal{P} := \{\tau_g : g \in \text{Sym}(n)\}$$

acting on  $\mathcal{G}_n$  where  $\tau_g(X) = X^g$ . The set of orbits is precisely  $\mathcal{H}$ . By Burnside's lemma,

$$|\mathcal{H}| = \frac{1}{n!} \sum_{\tau_g \in \mathcal{P}} |\text{fix}(\tau_g)|.$$

Consider some  $g \in \text{Sym}(n)$ .  $\text{fix}(\tau_g)$  is the set of graphs in  $\mathcal{G}_n$  fixed by  $\tau_g$ . Note that  $g$  induces a permutation  $\sigma_g$  on  $E(K_n)$ . Let  $\mathcal{C}$  be an orbit of the permutation  $\sigma_g$  and suppose  $X$  is fixed by  $\tau_g$ . Either  $X$  contains all edges in  $\mathcal{C}$  or has none of the edges in  $\mathcal{C}$ , otherwise, some edge will be mapped to a none-edge. Hence  $|\text{fix}(\tau_g)| = 2^{\text{orb}_2(g)}$  where  $\text{orb}_2(g)$  is the number of orbits in  $\sigma_g$ .

Among all permutations  $g \in \text{Sym}(n)$  whose support is size  $2r$ , the maximum value of  $\text{orb}_2(g)$  is realized by some permutation with  $r$  cycles of length 2. Suppose  $g$  is such a maximizer. Since  $g^2 = \text{id}$ , its action on  $E(K_n)$  has orbits of size at most 2. There are two ways in which an edge  $xy \in E(K_n)$  is NOT fixed by  $g$ . Either  $x, y$  are in the support of  $g$ , but in different orbits, or  $x$  is in the support of  $g$  and  $y$  is not. Hence the number of

orbits of length 2 in its action on  $E(K_v)$  is

$$r(r-1) - r(n-2r) = r(n-r-1)$$

and

$$\text{orb}_2(g) = \binom{n}{2} - r(n-r-1).$$

We now argue that the main contribution of Burnside's average comes from the identity permutation. Indeed,  $\text{orb}_2(\text{id}) = \binom{n}{2}$ , thus  $\sum_{\tau_g \in \mathcal{P}} |\text{fix}(\tau_g)| \geq 2 \binom{n}{2}$ .

Fix an even integer  $m \leq n-2$ . Consider the nonidentity permutations with support of size at most  $m$ . There are at most  $\binom{n}{m} m! < n^m$  permutations of this form. Let  $g$  be such an element. Then  $\text{orb}_2(g)$  is maximized if it has a single 2-cycle. It follows that

$$\text{orb}_2(g) \leq \binom{n}{2} - (n-2).$$

Of the rest of the permutations, there are at most  $n! < n^n$  of them. Let  $g$  be such a permutation.  $\text{orb}_2(g)$  is maximized at some permutation with  $\frac{m}{2}$  2-cycles, so

$$\text{orb}_2(g) \leq \binom{n}{2} - \frac{m}{2} \left( n - \frac{m}{2} - 1 \right) \leq \binom{n}{2} - \frac{nm}{4}$$

by our work above.

Therefore,

$$\begin{aligned} \sum_{\tau_g \in \mathcal{P}} |\text{fix}(\tau_g)| &\leq 2 \binom{n}{2} + n^m 2^{\binom{n}{2} - (n-2)} + n^n 2^{\binom{n}{2} - \frac{nm}{4}} \\ &= 2 \binom{n}{2} \left( 1 + n^m 2^{-(n-2)} + n^n 2^{-\frac{nm}{4}} \right). \end{aligned}$$

By taking  $m = \lfloor c \log n \rfloor$  for  $c > 4$ , we have

$$\begin{aligned} n^m 2^{-(n-2)} + n^n 2^{-\frac{nm}{4}} &= 2^{m \log n - n + 2} + 2^{n \log n - \frac{nm}{4}} \\ &= 2^{c \log^2 n - n + 2} + 2^{n \log n - \frac{cn \log n}{4}} \\ &= o(1). \end{aligned}$$

### Theorem 2.3.3

Asymptotically, almost all graphs are asymmetric.

That is, if we uniformly pick a graph in  $\mathcal{G}_n$  at random, the probability that the chosen graph is asymmetric tends to 1 as  $n \rightarrow \infty$ .

#### Proof

Let  $\mathcal{H}$  be the set of isomorphism classes of graphs on  $[n]$ . Pick  $\mathcal{C} \in \mathcal{H}$ . If graphs in  $\mathcal{C}$  are asymmetric, then  $|\mathcal{C}| = n!$ . Otherwise,  $|\mathcal{C}| \leq \frac{n!}{2}$ .

Let  $\rho$  be the proportion of all classes in  $\mathcal{H}$  such that  $|\mathcal{C}| = n!$ . Now,

$$\begin{aligned} 2^{\binom{n}{2}} &= |\mathcal{G}_n| \\ &= \sum_{\mathcal{C} \in \mathcal{H}} |\mathcal{C}| \\ &\leq \rho |\mathcal{H}| \cdot n! + (1 - \rho) |\mathcal{H}| \cdot \frac{n!}{2} \end{aligned}$$

By the previous lemma,

$$\begin{aligned} 2^{\binom{n}{2}} &\leq (1 + o(1)) \frac{2^{\binom{n}{2}}}{n!} \cdot n! \left( \rho + \frac{1 - \rho}{2} \right) \\ &= (1 + o(1)) 2^{\binom{n}{2}} \cdot \frac{1 + \rho}{2}. \end{aligned}$$

It follows that  $\rho \rightarrow 1$  as  $n \rightarrow \infty$  and the proportion of asymmetric graphs in  $\mathcal{G}_n$  is

$$\frac{\text{size of isomorphism class} \cdot \text{number of asymmetric isomorphism classes}}{|\mathcal{G}_n|}$$

$$\begin{aligned} &= \frac{n! \cdot \rho |\mathcal{H}|}{2^{\binom{n}{2}}} \\ &= (1 + o(1)) \rho \\ &\rightarrow 1. \end{aligned}$$

previous lemma

## 2.4 Primitivity

### Definition 2.4.1 (Block of Imprimitivity)

Let  $G$  be a group acting transitively on  $V$  and  $S \subseteq V$ .  $S$  is called a block of imprimitivity for  $G$  if  $\forall g \in G$ , either  $S^g = S$  or  $S^g \cap S = \emptyset$ .

### Example 2.4.1 (Trivial Blocks of Imprimitivity)

$S = \{u\}$  for some  $u \in V$  and  $S = V$  are trivial blocks of imprimitivity.

### Definition 2.4.2 (Primitive)

$G$  is primitive if there is no non-trivial blocks of imprimitivity.

If  $G$  is not primitive, we say it is *imprimitive*.

### Example 2.4.2

$\text{Aut}(K_n)$  has no non-trivial blocks of imprimitivity.

### Example 2.4.3

Consider  $Q_3$ , the 3-cube. Take any pair of vertices on  $S = \{a, b\}$  on opposite corners of a diagonal. Since automorphisms preserve distances,  $S$  is a block of imprimitivity.

Consider  $\text{Aut}(C_n)$  for composite  $n$ . Suppose  $n = xy$  for  $2 \leq x \leq y \in \mathbb{Z}_+$ . Consider  $S = \{0, x, 2x, \dots, (y-1)x\} \subseteq V$ . Since  $\text{Aut}(C_n)$  is precisely the dihedral group, any  $g \in \text{Aut}(C_n)$  is uniquely determined by its action on  $0, x$ . Thus  $S$  is a block of imprimitivity.

For  $n$  prime, we claim  $C_n$  is primitive. Let  $S \subsetneq V, |S| \geq 2$  and fix  $v \neq w \in S$  minimizing  $d_G(v, w)$  for all distinct pairs in  $S$ . Let  $r \in \text{Aut}(C_n)$  be the rotation mapping  $v \mapsto w$  (along the shortest  $vw$  path in  $G$ ). Clearly  $v^r = w$ . We claim that  $S^g \neq S$ . Otherwise, walking from  $w$  in the opposite direction to that of  $v$ , the first vertex of  $S$  we see is of distance  $d_C(v, w)$ . Following this argument, we see that  $S$  consists of consecutively equidistance vertices. But then  $n$  could not have been prime.

For groups  $H, K$ , we write  $H \leq K$  if  $H$  is a subgroup of  $K$  and  $H < K$  if  $H$  is a proper subgroup of  $K$ .

### Definition 2.4.3 (Maximal Subgroup)

$H < G$  is a maximal subgroup if there is no subgroup  $K$  such that  $H < K$ .

**Lemma 2.4.4**

Let  $G$  be a transitive permutation group on  $V$  and  $x \in V$ .  $G$  is primitive if and only if  $G_x$  is a maximal subgroup of  $G$ .

**Proof**

We prove the contrapositives.

( $\neg \implies \neg$ ) Let  $B$  be a block of imprimitivity where  $2 \leq |B| < |V|$ . By transitivity, there is some  $g \in G$  such that  $x \in B^g$ . Thus by taking  $B^g$  if necessary, we may assume without loss of generality that  $x \in B$ .

We wish to show that there is a subgroup  $H$  where  $G_x < H < G$ . Suppose  $G_B$  is the subgroup of all  $g \in G$  such that  $B^g = B$ . We argue that  $G_x < G_B < G$ .

Let  $g \in G_x$ . Then  $x \in B \cap B^g$ , hence  $B = B^g$  and  $G_x \leq G_B$ . But the fact that  $|B| \geq 2$  implies there is some  $y \neq x$  such that  $y \in B$ .

Let  $h \in G$  be such that  $x^h = y$ . This shows that  $h \notin G_x$ . Then  $y \in B \cap B^h$  and  $B = B^h$ . Thus  $h \in G_B$ . This shows that  $G_x < G_B$ .

The fact that  $G_B < G$  is obvious since  $B \subsetneq V$  and  $G$  is transitive.

( $\neg \longleftarrow \neg$ ) Suppose  $H$  is a subgroup of  $G$  such that  $G_x < H < G$ . Let  $B$  be the orbit of  $H$  containing  $x$ . We argue that  $B$  is a block of imprimitivity.

Now,  $G_x < H$  so there is some  $g \in H$  such that  $x^g \neq x$ . But  $x^g \in B^g$  then implies  $|B| \geq 2$ . On the other hand, the orbit-stabilizer formula states that

$$\begin{aligned}
 |B| &= \frac{|H|}{|H_x|} \\
 &< \frac{|G|}{|H_x|} \cdot \frac{|H_x|}{|G_x|} && H < G, H_x \subseteq G_x \\
 &= \frac{|G|}{|G_x|} \\
 &= |x^G| \\
 &= |V|. && \text{transitivity}
 \end{aligned}$$

Fix  $g \in G$ . We claim that  $B \cap B^g \neq \emptyset$  implies  $g \in H$ . Observe that  $B$  is  $H$ -invariant by the definition of an orbit. Thus the claim would show that  $B \cap B^g \neq \emptyset \implies B = B^g$ , concluding the proof.

Suppose  $y \in B \cap B^g$ . There is some  $h \in H$  such that  $x^h = y$ . Moreover, there is some

|  $h' \in H$  such that  $y = x^{gh'}$ . It follows that  $x = x^{h^{-1}gh'}$ . But then  $h^{-1}gh' \in G_x < H$  and  $g \in H$  as required.





# Chapter 3

## Transitive Graphs

### 3.1 Vertex-Transitive Graphs

**Definition 3.1.1 (Vertex Transitive)**

A graph  $X$  is vertex transitive if  $\text{Aut}(X)$  acts transitively on  $V(X)$ .

Let  $H(\cdot, \cdot)$  denote the Hamming distance between two strings in  $\mathbb{Z}_2^k$ .

**Definition 3.1.2 ( $k$ -Cube)**

The  $k$ -cube is the graph with vertex set

$$V(Q_k) := \mathbb{Z}_2^k$$

and edge set

$$E(Q_k) := \{\alpha\beta : H(\alpha, \beta) = 1\}.$$

We write  $Q_k$  to denote the  $k$ -cube.

**Lemma 3.1.1**

$Q_k$  is vertex transitive for all  $k \geq 1$ .

**Proof**

Fix  $v \in Q_k$ . For each  $x \in V(Q_k)$ , let  $P_v : V(Q_k) \rightarrow V(Q_k)$  be given by

$$x \mapsto x + v.$$

Observe that  $H(x, y) = 1 \iff H(x + v, y + v) = 1$ . Write  $P$  to denote the subgroup of automorphisms of the form  $P_v$ .

It follows that  $\text{Aut}(X) \geq P$  acts transitively on  $V(Q_k)$  by definition since the map  $P_{w-v}$  maps  $v \mapsto w$  for any  $v, w \in V(Q_k)$ .

As a corollary to this lemma, the  $|\text{Aut}(Q_k)| \geq 2^k$ . We improve upon this lowerbound.

**Proposition 3.1.2**

$|\text{Aut}(Q_k)| \geq 2^k k!$ .

**Proof**

Pick  $g \in \text{Sym}(k)$  and consider  $\tau_g : V(Q_k) \rightarrow V(Q_k)$  which permutes the entries of  $x \in V(Q_k)$  according to  $g$ .

We claim that

$$P' := \{\tau_g : g \in \text{Sym}(k)\} \leq \text{Aut}(Q_k).$$

To see this, observe that  $H(x, y) = 1 \iff H(\tau_g(x), \tau_g(y)) = 1$  for all  $x, y \in \mathbb{Z}_2^k$ .

Moreover, we claim that  $P' \cap P = \{\text{id}\}$ . This is because the ratio of 0-1 entries of a binary string  $x$  always remains the same under  $P'$ , but for all  $P_v \in P$ , we can find some string  $x$  for which the 0-1 ratio of  $x^{P_v}$  differs from  $x$ .

It follows by elementary group theory that

$$|PP'| = \frac{|P| \cdot |P'|}{|P \cap P'|} = 2^k k!$$

as desired.

Applying a similar proof, we can show that circulant graphs are also vertex transitive by the rotation subgroup to its automorphism group.

**Definition 3.1.3 (Cayley Graphs)**

Let  $G$  be a group and  $C \subseteq G$  which is closed under inverses and does NOT contain  $\text{id}$ .

The Cayley graph  $X = X(G, C)$  is defined by

$$\begin{aligned} V(X) &:= G \\ E(X) &:= \{gh : g^{-1}h \in C\}. \end{aligned}$$

Since  $\text{id} \notin C$ , Cayley graphs are loopless. If  $g^{-1}h \in C$ , then  $h^{-1}g = (g^{-1}h)^{-1} \in C$  as well.

Thus Cayley graphs are undirected.

**Theorem 3.1.3**

$X = X(G, C)$  is vertex transitive.

**Proof**

For each  $g \in G$ , let  $\tau_g : V(X) \rightarrow V(X)$  be given by

$$x \mapsto gx.$$

Observe that  $\tau_g \in \text{Aut}(X)$  as

$$\begin{aligned} xy \in E(X) &\iff y^{-1}x \in C \\ &\iff (y^{-1}g^{-1})gx \in C \\ &\iff \tau_g(y)^{-1}\tau_g(x) \in C \\ &\iff \tau_g(x)\tau_g(y) \in E(X). \end{aligned}$$

It follows that the set  $\{\tau_g : g \in G\}$  acts transitively on  $V(X)$ .

A special case of this is the  $k$ -cube where  $Q_k = X(\mathbb{Z}_2^k, \{e_i : i \in [k]\})$ .

A circulant graph  $X(\mathbb{Z}_n, C)$  is a Cayley graph under the same parameters.

**Proposition 3.1.4**

For  $v \geq k \geq i$ , the Johnson graph  $J(v, k, i)$  is vertex transitive.

**Proof**

For each  $g \in \text{Sym}(v)$ , let  $\sigma_g : V(J(v, k, i)) \rightarrow V(J(v, k, i))$  be given by

$$S \mapsto S^g.$$

Consider the subgroup

$$\{\sigma_g : g \in \text{Sym}(v)\} \leq \text{Aut}(J(v, k, i)).$$

By observation, the subgroup acts transitively on  $V(J(v, k, i))$ .

In general,  $J(v, k, i)$  are NOT Cayley graphs.

**Lemma 3.1.5**

The Peterson graph  $J(5, 2, 0)$  is NOT a Cayley graph.

**Proof**

There are only two groups of order 10, the cyclic group  $\mathbb{Z}_{10}$  and the dihedral group  $D_{10}$ .

Neither of the cubic Cayley graphs on these groups are isomorphic to  $J(5, 2, 0)$ .

Hence Cayley graphs are a strict subset of vertex transitive graphs.

## 3.2 Edge-Transitive Graphs

Recall that if  $G$  acts on  $V(X)$ , this induces an action on  $E(X)$ .

**Definition 3.2.1 (Edge Transitive)**

A graph  $X$  is edge transitive if  $\text{Aut}(X)$  acts transitively on  $E(X)$ .

If  $G$  acts on  $V(X)$ , this also induces an action on  $V(X)^2$ !

**Definition 3.2.2 (Arc Transitive)**

A graph  $X$  is arc transitive if  $\text{Aut}(X)$  acts transitively on  $\{(x, y) : xy \in E(X)\}$ .

Note that arc transitivity implies both vertex and edge transitivity.

In general, there is no relationship between vertex and edge transitive graphs.

**Example 3.2.1**

The triangular prism is vertex transitive but not edge transitive.

**Example 3.2.2**

For  $n \neq m$ ,  $K_{n,m}$  is edge transitive but not vertex transitive.

There are also examples of regular edge transitive graphs that are not vertex transitive.

**Lemma 3.2.3**

Suppose  $X$  is edge transitive and has no isolated vertices. If  $X$  is NOT vertex transitive,  $\text{Aut}(X)$  has exactly 2 orbits and they form a bipartition of  $X$ .

**Proof**

Aut( $X$ ) has exactly 2 orbits: First observe that  $\text{Aut}(X)$  has at least 2 orbits since it is not vertex transitive.

Fix  $xy \in E(X)$ . For any  $w \in V(X)$ , there is some  $wz \in E(X)$  since  $w$  is not isolated. Hence there is some  $\sigma \in \text{Aut}(X)$  such that

$$\{\sigma(x), \sigma(y)\} = \{w, z\}.$$

It follows that  $X \subseteq x^G \cup y^G$  by the choice of  $w$  and so  $X$  has at most 2 orbits.

The orbits yield a bipartition: We argue that there are no edges between vertices of  $x^G$  nor  $y^G$ . Suppose towards a contradiction that there is some  $wz \in E(X)$  where either  $w, z \in x^G$  or  $w, z \in y^G$ .

By the definition of orbits, there is no element of  $\text{Aut}(X)$  which maps  $wz \mapsto xy$ , as the latter has elements from both orbits. This is a contradiction since we assumed that  $X$  is edge transitive.

By contradiction,  $(x^G, y^G)$  forms a bipartition of  $X$ .

In general, edge transitivity and vertex transitivity does not imply arc transitivity.

**Definition 3.2.3 (Oriented Graph)**

An oriented graph is a directed graph  $X$  where  $xy \in A(X) \implies yx \notin A(X)$ .

The definition of automorphisms of undirected graphs naturally generalises to automorphisms of directed and therefore oriented graphs.

**Lemma 3.2.4**

If  $X$  is both vertex and edge transitive but not arc transitive, then the degrees of  $X$  are even.

**Proof**

For an undirected graph  $X$ , let  $D(X)$  be the directed graph obtained from  $X$  by replacing each  $uv \in E(X)$  with the two arcs  $(u, v), (v, u) \in A(D(X))$ .

Fix  $x \in V(X)$  and  $y \in N_X(x)$ . Put

$$\Omega_1 := (x, y)^G, \Omega_2 := (y, x)^G$$

where  $G = \text{Aut}(X)$ . Since  $X$  is edge transitive, then for any  $uv \in E(X)$ , either  $(u, v) \in \Omega_1$  or  $(v, u) \in \Omega_2$ . Hence  $\Omega_1 \cup \Omega_2 = A(D(X))$ .

But  $X$  is not arc transitive, thus  $\Omega_1 \cap \Omega_2 = \emptyset$ . It follows that  $(y, x) \notin \Omega_1$ . By the arbitrary choice of  $xy \in E(X)$ ,

$$(x, y) \in \Omega_1 \iff (y, x) \in \Omega_2.$$

Thus  $\Omega_1, \Omega_2$  are oriented graphs.

By construction,  $\text{Aut}(X) \leq \text{Aut}(\Omega_1)$  since every automorphism of  $X$  yields an automorphism of  $\Omega_1$ . Moreover,  $\text{Aut}(X)$  acts transitively on  $V(\Omega_1) = V(X)$  by assumption.

By transitivity,  $\Omega_1$  is necessarily a regular oriented graph. Thus for a vertex  $u \in V(X)$ .

$$\begin{aligned} \deg_X(u) &= d_{\Omega_1}^+(u) + d_{\Omega_1}^-(u) \\ &= 2d_{\Omega_1}^+(u) \\ &\equiv 0 \pmod{2}. \end{aligned}$$

#### Corollary 3.2.4.1

If  $X$  is vertex and edge transitive and the degrees of  $X$  are odd, then  $X$  is arc transitive.

### 3.3 Edge & Vertex Connectivity

#### Definition 3.3.1 (Edge Connectivity)

The edge connectivity of a graph  $X$ , denoted  $\kappa_1(X)$ , is the minimum cardinality of an edge subset whose deletion increases the number of components in  $X$ .

Let  $A \subseteq V(X)$ . Write

$$\partial A = \delta(A)$$

to be the set of edges leaving  $A$ .

#### Definition 3.3.2 (Edge Atom)

An edge atom of  $X$  is a minimum cardinality subset  $S \subseteq V(X)$  such that  $|\partial(S)| = \kappa_1(X)$ .

**Lemma 3.3.1**

Let  $X$  be connected and vertex transitive. Any two distinct edge atoms are vertex disjoint.

**Proof**

Let  $\kappa = \kappa_1(X)$ . Suppose  $A \neq B$  are edge atoms. By minimality,  $|A|, |B| \leq \frac{|V(X)|}{2}$  or else taking  $\bar{A}$  or  $\bar{B}$  contradicts minimality.

Case I:  $A \cup B = V(X)$  Then  $|A| = |B| = \frac{|V(X)|}{2}$  and we necessarily have  $A \cap B = \emptyset$ .

Case II:  $A \cup B \subsetneq V(X)$  Since  $A \subsetneq A \cup B$ , it must be that  $|\partial(A \cup B)| \geq \kappa$ . On the other hand,  $A \cap B \subsetneq A$  implies  $|\partial(A \cap B)| > \kappa$ .

By the submodularity of the  $|\partial(\cdot)|$  function

$$2\kappa < |\partial(A \cup B)| + |\partial(A \cap B)| \leq |\partial A| + |\partial B| = 2\kappa,$$

which is a contradiction.

**Lemma 3.3.2**

Let  $X$  be connected and vertex transitive. Suppose  $S$  is a block of imprimitivity for  $\text{Aut}(X)$ . Then  $X[S]$  is regular.

**Proof**

Let  $u \neq v \in S$  and set  $Y := X[S]$ . Since  $X$  is vertex transitive, we can choose some  $g \in \text{Aut}(X)$  for which  $u^g = v$  so that  $S \cap S^g = \emptyset$ .

By the choice of  $S$ ,  $S = S^g$ . Hence  $N_Y(v) = N_Y(u)^g$  and  $\deg_Y(u) = \deg_Y(v)$  as required.

**Theorem 3.3.3**

If  $X$  is connected and vertex-transitive, then  $\kappa_1(X)$  is equal to the degree of vertices in  $X$ .

**Proof**

Let  $k$  be the degrees of vertices in  $X$ . We must have  $\kappa_1(X) \leq k$ . We wish to argue that  $\kappa_1(X) \geq k$ . Let  $A$  be an edge atom. It is sufficient to show that  $\kappa_1(X) = |\partial A| \geq k$ .

Let  $g \in \text{Aut}(X)$  and  $B := A^g$ . We have shown that either  $A = B$  or  $A \cap B = \emptyset$  since  $B$  is also an edge atom. It follows that  $A$  is a block of imprimitivity.

But then  $X[A]$  is  $\ell$ -regular for some  $0 \leq \ell \leq k - 1$ . We know that  $\ell < k$  since  $X$  is connected. Hence  $|A| \geq \ell + 1$  and  $\ell \leq |A| - 1$ . Thus

$$\begin{aligned} |\partial A| &= |A|(k - \ell) && \ell \text{ neighbours within } A \\ &\geq |A|(k + 1 - |A|) \\ &\geq 1 \end{aligned}$$

and  $k \geq |A|$ .

Now then,

$$\begin{aligned} |\partial A| - k &\geq |A|k + |A| - |A|^2 - k \\ &= k(|A| - 1) - |A|(|A| - 1) \\ &= (|A| - 1)(k - |A|) \\ &\geq 0 \cdot 0 \\ &= 0. \end{aligned}$$

Notice we have equality if  $|A| = k$  or  $|A| = 1$ .

We conclude that  $\kappa_1(X) = k$  as required.

### Theorem 3.3.4

A  $k$ -regular vertex-transitive graph has vertex connectivity at least

$$\frac{2}{3}(k + 1).$$

We can find graphs for which this bound is tight.

## 3.4 Matchings

We say a vertex is *critical* if it is saturated in every maximum matching.

### Lemma 3.4.1

Let  $z_1, z_2 \in V(X)$  be vertices such that no maximum matching misses both of them. Suppose  $M_{z_1}, M_{z_2}$  are maximum matchings missing  $z_1, z_2$  respectively. Then  $M_{z_1} \Delta M_{z_2}$  has an even alternating  $z_1 z_2$ -path.



**Proof**

Let  $M_{z_1}, M_{z_2}$  be maximum matchings. Thus  $z_1, z_2$  are each on SOME even alternating path in  $M_{z_1} \Delta M_{z_2}$ .

Let  $P$  be the path containing  $z_1$ . If it does not contain  $z_2$ , then  $M_{z_2} \Delta P$  is a maximum matching missing both  $z_1, z_2$ . which is a contradiction.

**Lemma 3.4.2**

Let  $u, v \in V(X)$  be vertices in a connected graph  $X$  with no critical vertices. Then  $X$  contains a matching missing at most 1 vertex.

**Proof**

To see this we show that for every  $u, v \in V(X)$ , there is no maximum matching missing both  $u, v$ . Then any matching misses at most 1 vertex.

Let  $P$  be a  $uv$ -path. We argue by induction on the length  $|P|$ . The case where  $|P| = 1$  is trivial.

Suppose  $|P| \geq 2$ . Let  $x \notin \{u, v\}$  be a vertex on  $P$ . By the induction hypothesis, no maximum matching misses both  $u, x$  nor does a maximum matching miss both  $v, x$ .

Since  $x$  is not critical, there is a maximum matching  $M_x$  missing  $x$ . This matching must then saturate  $u, v$ .

Suppose now towards a contradiction that some maximum matching  $M$  misses both  $u, v$ . By the previous lemma,  $M_x \Delta M$  has an even alternating path  $ux$ -path. Similarly,  $M_x \Delta M$  has an even alternating  $xv$ -path.

By the parity of the paths and the fact that  $u, v$  are not  $M$  saturated, this can only happen if  $u = v$ , which contradicts our assumptions.

By induction, we conclude the result.

**Proposition 3.4.3**

Let  $X$  be connected and vertex transitive. Then  $X$  has a matching missing at most 1 vertex.

**Proof**

Case I:  $X$  has a critical vertex Let  $u$  be a critical vertex and any maximum matching  $M$ . Fix any other vertex  $v$ . Since  $X$  is vertex transitive, there is some  $g \in \text{Aut}(X)$  for which  $v^g = u$ .

$M^g$  is a maximum matching and thus saturates  $u$ . But then  $M$  saturates  $v$ .

Thus every vertex is critical and  $X$  has a perfect matching.

Case II:  $X$  has no critical vertices By our previous lemma,  $X$  contains a matching missing at most 1 vertex.

#### **Lemma 3.4.4**

Let  $X$  be connected and vertex transitive. Then every edge lies in some maximum matching.

#### **Proof**

We argue by induction on the number of vertices and the number of edges in a connected vertex-transitive graph. The base case of a single isolated vertex is trivial.

Inductively, suppose that some  $e \in E(X)$  is not contained in any maximum matching. It follows that  $X$  is not edge transitive. Let  $Y$  be the subgraph induced by the orbit of  $e$  under  $\text{Aut}(X)$ . Then  $Y$  must be a proper subgraph of  $X$  and we claim  $Y$  is vertex as well as edge transitive.

First,  $\text{Aut}(X) \leq \text{Aut}(Y)$ . Indeed, for all  $uv \in E(Y)$  and  $g \in \text{Aut}(X)$ ,

$$u^g v^g \in e^{\text{Aut}(X)} \implies u^g v^g \in E(Y).$$

But then  $Y$  is edge transitive by construction. Moreover, the vertex transitivity is due to the vertex transitivity of  $X$ .

Case I:  $Y$  is connected  $Y$  has a maximum matching missing at most 1 vertex. But then by the edge transitivity of  $Y$ , there is a maximum matching missing at most 1 vertex of  $V(Y) = V(X)$  that includes  $e$ . This is necessarily a maximum matching of  $X$  covering  $e$ , which is a contradiction.

Case II:  $Y$  has components  $C_1, C_2, \dots, C_m$  We claim that each  $C_i$  is a block of imprimitivity of  $\text{Aut}(X)$ . Moreover, the components are pairwise isomorphic and are vertex plus edge transitive.

The fact that each component is a block of imprimitivity follows from the fact that automorphisms preserve distances. Hence if a vertex within a component is sent to the same component, all vertices of that component are sent there as well.

The vertex and edge transitivity follows from each component being blocks of imprimitivity as well as the fact that  $Y$  is vertex and edge transitive. Finally, vertex transitivity and the assumption that each component is a block of imprimitivity yields the pairwise isomorphisms.

If each component has an even number of vertices, then each  $C_i$  has a perfect matching. But then by edge transitivity,  $Y$  and thus  $X$  has a perfect matching containing  $e$ .

Suppose now that each component is odd. Define

$$\begin{aligned} V(Z) &= \{C_1, \dots, C_m\} \\ E(Z) &= \{C_i C_j : \exists f \in \partial_X(C_i) \cap \partial_X(C_j)\} \end{aligned}$$

$Z$  must be connected since  $X$  is connected. We claim that  $Z$  is vertex transitive.

To see this claim recall how  $\text{Aut}(X)$  permutes the components. Thus  $f \in \partial_X(C_i) \cap \partial_X(C_j)$ , implies  $f^g \in \partial_X(C_i^g) \cap \partial_X(C_j^g)$  for all  $g \in \text{Aut}(X)$ . Hence  $\text{Aut}(X)$  induces a subgroup of  $\text{Aut}(Z)$  which acts transitively on the vertices.

Now,  $Z$  has a matching missing at most 1  $C_i$ . By vertex transitivity, we may assume this is the component  $C$  containing the edge  $e$ . By the edge transitivity of  $C$ , there is a maximum matching containing  $e$  and missing at most 1 vertex.

For each  $C_i C_j$  in the matching, let us pick an edge in  $\partial_X(C_i) \cap \partial_X(C_j)$ . This is a matching  $M$  in  $X$ . Using the fact that each  $C_i$  is vertex transitive with a maximum matching missing 1 vertex, we can thus use  $M$  to “connect” the maximum matchings within  $C_i$ . This yields a maximum matching of  $X$  missing exactly 1 vertex and containing  $e$ .

By induction, we conclude that every edge is contained in some maximum matching.

### Theorem 3.4.5

Let  $X$  be connected and vertex transitive. Then  $X$  has a matching missing at most 1 vertex and each edge is contained in some maximum matching.

## 3.5 More on Cayley Graphs

### Definition 3.5.1 (Semiregular)

A permutation group  $G$  acting on  $V$  is semiregular if  $G_x = 1$  for all  $x \in V$ .

### Proposition 3.5.1

If  $G$  is semiregular, then

$$|G| = |x^G|$$

for all  $x \in V$ .

**Proof**

Orbit-Stabilizer lemma.

**Definition 3.5.2 (Regular)**

If  $G$  is semiregular and transitive, then  $G$  is regular.

**Proposition 3.5.2**

If  $G$  acting on  $V$  is regular, then  $|G| = |V|$ .

**Proof**

$$|G| = |x^G| = |V|.$$

**Lemma 3.5.3**

Let  $G$  be a group and  $C \subseteq G \setminus \{1\}$  be inverse-closed. Then  $\text{Aut}(X(G, C))$  contains a regular subgroup isomorphic to  $G$ .

**Proof**

Let  $X = X(G, C)$ . For  $g \in G$ , define  $\tau_g : V(X) \rightarrow V(X)$  given by

$$x \mapsto gx.$$

Write  $T := \{\tau_g : g \in G\}$ . By our initial work with Cayley graphs, we know that  $T \leq \text{Aut}(X)$ . Moreover,  $T$  acts transitively on  $X$ . In addition  $T \cong G$  by the map  $\tau_g \mapsto g$ . Finally,  $T$  is semiregular, as any non-identity element does not fix any vertices.

Thus by definition,  $T$  is a regular subgroup of  $\text{Aut}(X)$ .

**Lemma 3.5.4**

If  $G \leq \text{Aut}(X)$  acts regularly on  $V(X)$ , then  $X = X(G, C)$  for some inverse-closed  $C \subseteq G \setminus \{1\}$ .

**Proof**

Since  $G$  is regular, then  $|G| = |V(X)|$ . We aim to construct  $C$  so that  $X \cong X(G, C)$ .

Fix  $u \in V(X)$ . Since  $G$  is regular, there is a unique  $g_v \in G$  such that  $u^{g_v} = v$ . Define

$$C := \{g_v : v \sim u\}$$

as the subset of  $G$  sending  $u$  to a neighbour of  $u$ .

It is clear that  $1 \notin C$ , as  $X$  does not have loops. It suffices to show that

$$xy \in E(X) \iff g_x^{-1}g_y \in C \iff g_y^{-1}g_x \in C.$$

$xy \in E(X) \iff g_x^{-1}g_y \in C$  Recall that  $g_x \in G \leq \text{Aut}(X)$ . Hence  $g_x^{-1} \in \text{Aut}(X)$ . Thus

$$x^{g_x^{-1}}y^{g_x^{-1}} \in E(X) \iff xy \in E(X).$$

Since  $u^{g_x} = x$ , we have  $x^{g_x^{-1}} = u$ . Similarly,  $u^{g_y} = y$  so  $y^{g_x^{-1}} = u^{g_x^{-1}g_y}$ . So

$$xy \in E(X) \iff uu^{g_x^{-1}g_y} \in E(X) \iff g_x^{-1}g_y \in C$$

as desired.

$xy \in E(X) \iff g_y^{-1}g_x \in C$  This is identical except we consider  $g_y$  instead of  $g_x$ .

Hence  $C$  is inverse closed and  $X \cong X(G, C)$  with an isomorphism given by  $x \mapsto g_x$ .

We are now concerned with Cayley graphs on the same group.

### Definition 3.5.3 (Group Automorphism)

Let  $G$  be a group.  $\theta : G \rightarrow G$  is an automorphism if  $\theta$  is bijective and  $\theta(gh)\theta(g)\theta(h)$  for all  $g, h \in G$ .

### Lemma 3.5.5

Let  $\theta$  be an automorphism of  $G$ . Then  $X(G, C) \cong X(G, \theta(C))$ .

#### Proof

We claim that  $\theta$  is an isomorphism  $X(G, C) \rightarrow X(G, \theta(C))$ .

We have

$$\begin{aligned} g^{-1}h \in C &\iff \theta(g^{-1}h) \in \theta(C) \\ &\iff \theta(g^{-1})\theta(h) \in \theta(C) \\ &\iff \theta(g)^{-1}\theta(h) \in \theta(C) \end{aligned}$$

as desired.

We will note that if  $X(G, C) \cong X(G, \vartheta(C))$ , then it is not necessary that  $\vartheta$  is an automorphism.

**Definition 3.5.4 (Generating Set)**

Let  $G$  be a group.  $C \subseteq G$  is a generating set if every  $g \in G$  is a product of elements in  $C$ .

**Lemma 3.5.6**

$X = X(G, C)$  is connected if and only if  $C$  is a generating set for  $G$ .

**Proof**

( $\implies$ ) If  $X$  is connected, there is a  $eg$ -path for all  $g \in G$ . Fix  $g \in G$ . Enumerate some  $eg$ -path

$$e = h_0, h_1, h_2, \dots, h_k, g.$$

Observe that  $h_{i-1}^{-1}h_i \in C$  for  $i \in [k]$ , and  $h_k^{-1}g \in C$ .

But then

$$g = \left( \prod_{i=1}^k h_{i-1}^{-1}h_i \right) h_k^{-1}g$$

is a product of elements of  $C$ .

( $\impliedby$ ) Suppose  $C$  generates  $G$ . We argue there is a  $eh$ -path for all  $h \in G$ .

Write  $h$  as a product of  $g_i$ 's from  $C$ :

$$h = c_1c_2 \dots c_k.$$

Notice that

$$\begin{array}{ll} e \sim c_1 & ec_1 = c_1 \in C \\ c_1 \sim c_1c_2 & c_1^{-1}c_1c_2 = c_2 \in C \\ \prod_{i=1}^{\ell} c_i \sim \prod_{i=1}^{\ell+1} c_i & \left( \prod_{i=1}^{\ell} c_i \right)^{-1} \prod_{i=1}^{\ell+1} c_i = c_{\ell+1} \in C \end{array}$$

This  $e$  is indeed connected to  $\prod_{i=1}^k c_i = h$  and we are done.

## 3.6 Retracts

### Definition 3.6.1 (Retract)

A subgraph  $Y$  of  $X$  is a retract if there is some homomorphism  $f : X \rightarrow Y$  such that

$$f \Big|_Y = 1_Y.$$

### Theorem 3.6.1

Every connected vertex-transitive graph is isomorphic to a retract of a Cayley graph.

### Proof

Let  $X$  be connected and vertex transitive. Fix  $x \in V(X)$  and

$$C := \{g \in \text{Aut}(X) : x \sim x^g\}.$$

Let  $G \leq \text{Aut}(X)$  as the subgroup generated by  $C$ . We claim that  $G$  acts transitively on  $V(X)$ .

It suffices to show that there is an element sending  $x \mapsto y$  for all  $y \in V(X)$ . We argue by induction on  $d(x, y)$  that  $y \in x^G$ . The case where  $d(x, y) = 1$  is trivial.

Suppose  $d(x, y) \geq 2$ . There is some neighbour of  $y$ , say  $z \neq x, y$ , on the shortest  $xy$ -path in  $X$ . By induction, there is some  $g \in G$  sending  $x \mapsto z$ . Since  $X$  is vertex transitive, we can find some  $h \in \text{Aut}(X)$  sending  $z \mapsto y$ .

We have  $z \sim z^h$ . Since automorphisms preserve edges,  $z^{g^{-1}} \sim z^{g^{-1}h}$  and  $g^{-1}h \in C \subseteq G$  by definition. But then  $h = g \cdot g^{-1}h \in G$  as both are elements of  $G$ .

Let  $Y := X(G, C)$ . We argue that  $X$  is isomorphic to a retract of  $Y$ .

Recall from our work with group actions on graphs that  $H := \{h \in G : x^h = y\}$  is a left coset of the stabilizer  $G_x$ . Thus for every  $y \in V(X)$ ,

$$C_y := \{g \in G : x^g = y\}$$

is a left coset of  $G_x$ . Hence  $C = \bigcup_{y \sim x} C_y$  by definition is a union of left cosets of  $G_x$ . Moreover,  $C \cap G_x = \emptyset$  as  $x \not\sim x$ .

Observe that for any  $a, b \in \text{Aut}(X)$ ,

$$x^a \sim x^b \iff x \sim x^{a^{-1}b} \iff a^{-1}b \in C.$$

Let  $A_1, \dots, A_k$  be the left cosets of  $G_x$ . Pick a representative  $a_i \in A_i$  for each  $i \in [k]$ . We claim that  $Y[a_1, \dots, a_k] \cong X$  and that  $Y[a_1, \dots, a_k]$  is a retract of  $Y$ . This would terminate the proof.

First, we show that  $C = G_x C G_x$ . Indeed, it is clear that  $C \subseteq G_x C G_x$ . Pick  $h, h' \in G_x$  and  $g \in C$ . We have

$$\begin{aligned} x &\sim x^g \\ x &= x^h \sim (x^h)^g = x^{gh} \\ x &= x^{h'} \sim x^{h'gh}. \end{aligned}$$

Then  $h'gh \in C$  and  $G_x C G_x \subseteq C$  as required.

Next, we claim that in  $Y = X(G, C)$ ,  $E(A_i) = \emptyset$  for all  $i$ . Moreover, for all  $1 \leq i < j \leq k$ , either  $\partial A_i \cap \partial A_j = \emptyset$  or  $A_i, A_j$  induces a complete bipartite graph.

For any  $g' \in G$ , there is some  $j$  for which  $g' \in A_j$ . Thus we can write  $g' = a_j g$  for some  $g \in G_x$ . Suppose  $g, h \in G_x$ . Then

$$\begin{aligned} a_i g \sim a_j h &\iff (a_i g)^{-1} a_j h \in C \\ &\iff g^{-1} a_i^{-1} a_j h \in C \\ &\iff a_i^{-1} a_j \in g C h^{-1} \in G_x C G_x = C. \end{aligned}$$

This shows that the adjacency between any vertices in different cosets depend on the cosets and not the specific vertices. Moreover, since  $1 \notin C$ , we know that  $a_i^{-1} a_i \notin C$  and hence  $a_i g \not\sim a_i h$  for any  $g, h \in G_x$ .

We now claim that  $Y[a_1, \dots, a_k] \sim X$ . By the previous claim,  $a_i a_j \in E(Y[a_1, \dots, a_k])$  if and only if  $a_i^{-1} a_j \in C$ . Consider the map  $\rho : V(X) \rightarrow \{a_1, \dots, a_k\}$  given by

$$y \mapsto a_j \in C_y.$$

Select  $u, v \in V(X)$ . Since  $G$  acts transitively, there are  $g, h \in G$  such that  $x^g = u, v^h = v$ .

$$\begin{aligned} u \sim v &\iff x^g \sim x^h \\ &\iff x \sim x^{g^{-1}h} \\ &\iff g^{-1}h \in C. \end{aligned}$$

On the other hand,  $\rho(u)\rho(v) \in E(Y[a_1, \dots, a_k])$  if and only if  $\rho(u)^{-1}\rho(v) \in C$ . Since  $u = x^g, v = x^h$ , we know that  $g, h$  are representatives of the cosets to which  $\rho(u), \rho(v)$  belong, respectively. Thus we may assume without loss of generality that  $h = \rho(u), h = \rho(v)$ . So  $\rho(u) \sim \rho(v) \iff g^{-1}h \in C$  and  $\rho$  preserves adjacencies.



Finally, we show that  $Y[a_1, \dots, a_k]$  is a retract of  $Y$ . This is given by the homomorphism  $\tau : V(Y) \rightarrow \{a_1, \dots, a_k\}$  given by

$$g \mapsto g_j, g \in A_j.$$

Clearly then this is a homomorphism by our work above and acts as the identity on  $\{a_1, \dots, a_k\}$ .

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# Chapter 4

## Generalized Polygons

### 4.1 Incidence Graphs

**Definition 4.1.1 (Incidence Structure)**

Let  $\mathcal{P}$  be a set of points,  $\mathcal{L}$  a set of lines, and  $I \subseteq \mathcal{P} \times \mathcal{L}$  an incidence relation. The triple  $\mathcal{I} = (\mathcal{P}, \mathcal{L}, I)$  defines an incidence structure.

If  $(p, L) \in I$ , then we say point  $p$  and line  $L$  are *incident*.

The triple  $\mathcal{I}^* = (\mathcal{L}, \mathcal{P}, I^*)$  where

$$I^* = \{(L, p) : (p, L) \in I\}$$

is the *dual* of  $\mathcal{I}$ .

**Definition 4.1.2 (Incidence Graph)**

Given an incidence structure  $\mathcal{I} = (\mathcal{P}, \mathcal{L}, I)$ , the incidence graph  $X(\mathcal{I})$  is defined as the bipartite graph  $(\mathcal{P}, \mathcal{L}, E)$  where

$$pL \in E \iff (p, L) \in I.$$

Remark that  $X(\mathcal{I}) \cong X(\mathcal{I}^*)$ .

**Definition 4.1.3 (Automorphism)**

An automorphism of  $(\mathcal{P}, \mathcal{L}, I)$  is a permutation of  $\mathcal{P}, \mathcal{L}$  such that  $\mathcal{P}^\sigma = \mathcal{P}$  and  $\mathcal{L}^\sigma = \mathcal{L}$  with

$$(p^\sigma, L^\sigma) \in I \iff (p, L) \in I.$$

## 4.1.1 Partial Linear Space

**Definition 4.1.4 (Partial Linear Space)**

$\mathcal{I} = (\mathcal{P}, \mathcal{L}, I)$  is a partial linear space if for all  $x, y \in \mathcal{P}$ , there exists at most one  $L \in \mathcal{L}$  such that  $(x, L) \in I, (y, L) \in I$ .

**Lemma 4.1.1**

If  $\mathcal{I}$  is a partial linear space, then any two lines are incident with at most 1 point.

If lines  $L_1, L_2$  are incident with  $p$ , then we say these two lines are *concurrent* and meet at  $p$ .

**Lemma 4.1.2**

If  $\mathcal{I}$  is a partial linear space, then  $X(\mathcal{I})$  has girth at least 6.

**Proof**

$X(\mathcal{I})$  has no 4-cycle as otherwise, there are two points incident with two lines.

## 4.2 Projective Planes

**Definition 4.2.1 (Projective Plane)**

A projective plane is a partial linear space satisfying

- (C1) Every two lines meet in a unique point
- (C2) Every two points lie in a unique line
- (C3) There exists 3 noncollinear points (they form a triangle)

**Theorem 4.2.1**

Let  $\mathcal{I}$  be a partial linear space containing a triangle. Then  $\mathcal{I}$  is a projective plane if and only if  $X(\mathcal{I})$  has diameter 3 and girth 6.

**Proof**

( $\implies$ ) Let  $\mathcal{I} = (\mathcal{P}, \mathcal{L}, I)$  be a projective plane with a triangle. Then every  $x, y \in \mathcal{P}$  are at distance 2 and every  $L_1, L_2 \in \mathcal{L}$  are at distance 2.

Pick  $x \in \mathcal{P}, L \in \mathcal{L}$ . Either  $(x, L) \in I$ , or there is some  $L' \neq L$  such that  $(x, L') \in I$ . In the latter case, there is  $y \in \mathcal{P}$  such that  $L, L'$  meet at  $y$ . So  $d(x, L) \leq 3$  and the diameter is at most 3. But since  $\mathcal{I}$  has a triangle,  $X(\mathcal{I})$  is NOT a complete bipartite graph. Thus  $X(\mathcal{I})$  has diameter at least 3.

Now,  $\mathcal{I}$  is a partial linear space. Hence we know the girth is at least 6.

Let  $\{x, y, z\}$  be a triangle. Suppose  $x, y$  meet at  $L_1$ ,  $x, z$  meet at  $L_2$ , and  $y, z$  meet at  $L_3$ . Since  $x, y, z$  are not collinear,  $L_1, L_2, L_3$  are distinct. Hence the girth is at most 6.

( $\impliedby$ ) Conversely, suppose  $\mathcal{I}$  is a partial linear space and  $X(\mathcal{I})$  has diameter 3 as well as girth 6.

To see (C2), note that for any  $x \neq y \in \mathcal{P}$ , the distance is even and at most 3. Hence they are distance 2 apart. It follows that there is  $L \in \mathcal{L}$  such that  $(x, L), (y, L) \in I$ . Moreover, the girth being 6 (strictly more than 4) implies that such  $L$  is unique.

(C1) is satisfied with an identical argument.

(C3) is given by the assumption of girth being 6.

### 4.3 A Family of Projective Planes

Consider  $\mathbb{F}_q$ , the finite field of order  $q$ . Let  $V := \mathbb{F}_q^3$ . Define

$$\text{PG}(2, q) := (\mathcal{P}, \mathcal{L}, I)$$

where  $\mathcal{P}$  are the 1-dimensional subspaces of  $V$ ,  $\mathcal{L}$  are the 2-dimensional subspaces of  $V$ , and  $(p, \mathcal{L}) \in I$  if  $p$  is a subspace of  $L$ .

Alternatively,  $L$  can be represented by the orthogonal complement of  $L_+$ , which is  $\text{span}\{a\}$ . Here  $a^T p_1, a^T p_2 = 0$ .

Now,  $V$  has  $q^3 - 1$  non-zero vectors. There are  $q - 1$  non-zero vectors in each 1-dimensional

subspace. Thus

$$|\mathcal{P}| = \frac{q^3 - 1}{q - 1} = 1 + q + q^2.$$

By the bijection between 1-dimensional subspaces and 2-dimensional subspaces,

$$|\mathcal{L}| = 1 + q + q^2.$$

Each line has  $q^2 - 1$  non-zero vectors. Hence each line is incident with  $\frac{q^2-1}{q-1} = 1 + q$  points. Moreover, for  $(p, L) \in I$  to hold, we have  $p \subseteq L$ . There are  $q^2 - 1$  linearly independent vectors to  $p$ , each spanning  $q - 1$  non-zero points. Hence each point has  $1 + q$  lines passing through it.

Observe that  $X(\text{PG}(2, q))$  has  $2(q^2 + q + 1)$  vertices and is  $(q + 1)$ -regular.

Remark that  $\text{PG}(2, 2)$  is the Fano plane.

**Lemma 4.3.1**

$\text{PG}(2, q)$  is a projective plane.

**Proof**

Let  $L_1 := \text{span}\{u, v\}$  and  $L_2 := \text{span}\{u', v'\}$  such that  $L_1 \neq L_2$ . We have  $\dim(L_1 \cap L_2) \geq 1$  since  $\dim V = 3$ . But  $\dim(L_1 \cap L_2) < 2$  since  $L_1 \neq L_2$ . It follows that  $\dim(L_1 \cap L_2) = 1$ . Thus there is a unique point where  $L_1, L_2$  meet. This verifies (C1).

Let  $p_1 = \text{span}\{u\}$  and  $p_2 = \text{span}\{v\}$ , where  $v \notin p_1$ . Suppose  $L$  is a line incident with both points. Then  $\text{span}\{u, v\} \subseteq L$  and the fact that  $\dim L = 2$  implies that  $L = \text{span}\{u, v\}$ . This shows (C2).

Let  $u, v, w$  be 3 linearly independent vectors. Then  $\text{span}\{u, v, w\} = V$  and so  $u, v, w$  are not contained in a 2-dimensional subspace. Hence the spans of each vector form a triangle. This demonstrates (C3).

### 4.3.1 Automorphisms

Consider  $\text{GL}(3, q)$ , the set of invertible  $3 \times 3$  matrices over  $\mathbb{F}_q$ . Note that  $\text{GL}(3, q)$  is a group which acts on vectors of  $\mathbb{F}_q^3$  through multiplication.

**Lemma 4.3.2**

Let  $A \in \text{GL}(3, q)$  and  $X = X(\text{PG}(2, q))$ . Then  $A$  induces a permutation on  $V(X)$ .

**Lemma 4.3.3**

$$\text{GL}(3, q) \leq \text{Aut}(\text{PG}(2, q)).$$
**Proof**

Let  $A \in \text{GL}(3, q)$ . Assume that  $p \sim L$ , ie  $p \subseteq L$ . Suppose that  $p = \text{span}\{u\}$  and  $L = \text{span}\{u, v\}$ . Thus

$$\begin{aligned} p^A &= \text{span}\{Au\} \\ L^A &= \text{span}\{Au, Av\} \\ p^A &\sim L^A. \end{aligned}$$

Finally,  $\mathcal{P}^A = \mathcal{P}$  and  $\mathcal{L}^A = \mathcal{L}$ .

**Lemma 4.3.4**

$$X(\text{PG}(2, q)) \text{ is arc-transitive.}$$

Note this implies the incidence graph is both vertex and edge-transitive.

**Proof**

Suppose  $p_1 \sim L$  and  $p_2 \sim L_2$ . Write

$$\begin{aligned} p_1 &= \text{span}\{u_1\} \\ L_1 &= \text{span}\{u_1, v_1\} \\ p_2 &= \text{span}\{u_2\} \\ L_2 &= \text{span}\{u_2, v_2\}. \end{aligned}$$

By elementary linear algebra, there is some invertible  $A$  for which  $Au_1 = u_2, Av_1 = v_2$ . Then  $(p_1, L_1)^A = (p_2, L_2)$ . Hence  $\text{GL}(3, q)$  acts transitively on  $\mathcal{P}, \mathcal{L}$ .

It suffices to show that there is some  $\pi \in \text{Aut}(X)$  such that  $\pi(\mathcal{P}) = \mathcal{L}$ . This is given by

$$\pi \text{span}\{u\} := \text{span}\{u\}^\perp.$$

Suppose  $(p, L) \in I$  so  $p = \text{span}\{u\}, L = \text{span}\{u, v\}$ . We have

$$p^\pi = \text{span}\{u\}^\perp, L^\pi = \text{span}\{u, v\}^\perp.$$

But then  $L^\pi \subseteq p^\pi$  hence  $L^\pi \sim p^\pi$  and  $\pi \in \text{Aut}(X)$ .

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# Chapter 5

## Homomorphisms

### 5.1 Definitions & Basic Results

We write  $X \rightarrow Y$  to denote the existence of a homomorphism from  $X$  to  $Y$ .

**Lemma 5.1.1**

The relation  $\rightarrow$  is reflexive and transitive.

But neither symmetric nor antisymmetric. Thus  $X \rightarrow Y$  does not in general imply  $Y \rightarrow X$  and it is possible for  $X \not\cong Y$  to satisfy  $X \rightarrow Y, Y \rightarrow X$ .

Since  $\rightarrow$  is not symmetric, it is NOT a equivalence relation. Since  $\rightarrow$  is not antisymmetric, it is not a partial order, unless we restrict to a specific family of graphs.

**Definition 5.1.1 (Homomorphically Equivalent)**

We say  $X, Y$  are homomorphically equivalent if  $X \rightarrow Y, Y \rightarrow X$ .

This clearly an equivalence relation.

A homomorphism is *surjective* if  $V(Y) = f(V(X))$ .

**Lemma 5.1.2**

If surjective homomorphisms exist from  $X$  to  $Y$  and  $Y$  to  $X$ , then  $X \cong Y$ .

**Definition 5.1.2 (Odd Girth)**

The odd girth of  $X$  is the length of the shortest odd cycle in  $X$ . If  $X$  is bipartite, then its odd girth is  $\infty$ .

**Lemma 5.1.3**

Suppose  $X \rightarrow Y$ . Then

- (a)  $\chi(X) \leq \chi(Y)$
- (b) the odd girth of  $X$  is at least the odd girth of  $Y$

**Corollary 5.1.3.1**

- (a) There is no homomorphism from  $C_{2n+1}$  to  $K_2$
- (b) There is no homomorphism from the Peterson graph to  $C_4$
- (c) There is no homomorphism from the Peterson graph to any of its proper subgraph

## 5.2 Cores

**Definition 5.2.1 (Core)**

A graph  $X$  is a core if any automorphism from  $X$  to itself is an automorphism.

**Definition 5.2.2 (Core Of)**

A subgraph  $Y$  of  $X$  is a core of  $X$  if  $Y$  is a core and  $X \rightarrow Y$ .

Thus the cores of  $X$  are minimal subgraphs that are homomorphic images of  $X$ .

**Example 5.2.1**

$K_n$  is a core.

**Example 5.2.2**

A graph is *critical* if any of its proper subgraphs has strictly smaller chromatic number.

Then  $\chi(X) > \chi(Y)$  for every proper subgraph  $Y$  and  $X \not\rightarrow Y$ . Thus every critical graph is a core.

**Lemma 5.2.3**

If  $Y$  is a core of  $X$ , then  $Y$  is a retract of  $X$ .

**Proof**

Let  $f : X \rightarrow Y$  be a homomorphism. Since  $Y$  is a core,  $g := f|_Y$  is an automorphism of  $Y$ .

Then  $g^{-1} \circ f : X \rightarrow Y$  is a retraction.

**Lemma 5.2.4**

Let  $Y_1, Y_2$  be cores. Then  $Y_1, Y_2$  are homomorphically equivalent if and only if they are isomorphisms.

**Proof**

Let  $f : Y_1 \rightarrow Y_2$  and  $g : Y_2 \rightarrow Y_1$  be homomorphisms. Then both  $f \circ g, g \circ f$  are homomorphisms. But  $Y_1, Y_2$  are cores. So  $f \circ g$  and  $g \circ f$  are both bijective. Hence  $f, g$  must both be surjective.

**Lemma 5.2.5**

Every graph  $X$  has a core which is an induced subgraph and unique up to isomorphism.

Due to this result, we may refer to THE core  $X^\bullet$  of  $X$ .

**Proof**

Existence: Let  $\mathcal{F}$  be the set of subgraphs of  $X$  that are homomorphic images of  $X$ .  $\mathcal{F} \neq \emptyset$  since  $X$  is a homomorphic image of  $X$ .

Let  $Y$  be a minimal graph in  $\mathcal{F}$  with respect to the subgraph relation. By the minimality, any homomorphism from  $Y$  to itself must be an automorphism of  $Y$ . Hence  $Y$  is a core of  $X$ .

Since  $Y$  is a retract of  $X$ ,  $Y$  must be induced.

Uniqueness: Let  $Y_1, Y_2$  be cores of  $X$ . Let  $f_i : X \rightarrow Y_i$  be a homomorphism from  $X$  to  $Y_i$  for  $i = 1, 2$ .

Consider  $f_1|_{Y_2} : V(Y_2) \rightarrow V(Y_1)$ . This is a homomorphism from  $Y_2$  to  $Y_1$ . Similarly,

$f_2|_{Y_1}$  is a homomorphism from  $Y_1$  to  $Y_2$ . Hence  $Y_1 \rightarrow Y_2$  and  $Y_2 \rightarrow Y_1$  and they are homomorphically equivalent.

Since  $Y_1, Y_2$  are cores, they are homomorphically equivalent if and only if they are isomorphic.

**Lemma 5.2.6**

Two graphs  $X, Y$  are homomorphically equivalent if and only if their cores are isomorphic.

**Proof**

Suppose  $X \rightarrow Y, Y \rightarrow X$ . Then

$$\begin{aligned} X^\bullet &\rightarrow X \rightarrow Y \rightarrow Y^\bullet \\ Y^\bullet &\rightarrow Y \rightarrow X \rightarrow X^\bullet \end{aligned}$$

By the a lemma,  $X^\bullet \cong Y^\bullet$ .

Suppose now that  $X^\bullet \cong Y^\bullet$ . Then  $X^\bullet \rightarrow Y^\bullet$  and  $Y^\bullet \rightarrow X^\bullet$ .

It follows by a previous lemma that

$$\begin{aligned} X^\bullet &\rightarrow X \rightarrow Y \rightarrow Y^\bullet \\ Y^\bullet &\rightarrow Y \rightarrow X \rightarrow X^\bullet \end{aligned}$$

**Corollary 5.2.6.1**

$\rightarrow$  is a partial order on the set of isomorphism classes of cores.

**Proof**

We know  $\rightarrow$  is reflexive and transitive. It is antisymmetric on the set of isomorphism classes of cores since if  $X, Y$  are cores that are homomorphically equivalent, then  $X \cong Y$ .

**Lemma 5.2.7**

Let  $X$  be connected. If every path of length 2 lies in a shortest odd cycle of  $X$ , then  $X$  is a core.

**Proof**

Suppose towards a contradiction that  $X^\bullet \neq X$  and  $f : X \rightarrow X^\bullet$  is a retraction. Then by connectedness, there are  $u, v \in V(X)$  such that  $u \sim v$ ,  $v \in V(X^\bullet)$ , and  $u \notin V(X^\bullet)$ .

Let  $w := f(u) \in V(X^\bullet)$ . Then  $w \sim v$  and  $w \not\sim u$ . By assumption  $uvw$  is contained in a shortest odd cycle  $C$ . The homomorphic image of  $C$  under  $f$  is an odd closed walk where  $f(u) = f(w) = w$ .

Hence  $f(X)$  has a strictly shorter odd cycle, which is a contradiction by earlier lemma. By contradiction,  $X^\bullet = X$ .

### Corollary 5.2.7.1

The Peterson graph is a core.

## 5.3 Products

### Definition 5.3.1 (Product Graph)

For graphs  $Y, Z$ , the (direct) product  $Y \times Z$  is defined by  $V(Y \times Z) = V(Y) \times V(Z)$  where

$$(y, z) \sim (y', z') \iff y \sim y' \wedge z \sim z'.$$

### Lemma 5.3.1

- (a) If  $Y, Z$  are both connected. Then  $Y \times Z$  are disconnected if and only if both  $Y, Z$  are bipartite.
- (b) Let  $Y_1 + Y_2$  denote the disjoint union of  $Y_1, Y_2$ . Then  $(Y_1 + Y_2) \times Z = Y_1 \times Z + Y_2 \times Z$ .
- (c)  $Y \times Z \cong Z \times Y$  (commutative).  $(Y_1 \times Y_2) \times Y_3 \cong Y_1 \times (Y_2 \times Y_3)$  (associative). However,  $Y \times Z_1 \cong Y \times Z_2$  does not imply  $Z_1 \cong Z_2$  [ie  $K_1 \times (K_3 + K_3) \cong K_2 \times C_6 \cong C_6 + C_6$ ].
- (d)  $p_X : V(X \times Y) \rightarrow V(X)$  and  $p_Y : V(X \times Y) \rightarrow V(Y)$  given by  $(x, y) \mapsto x, (x, y) \mapsto y$ , respectively, are homomorphisms from  $X \times Y$  to  $X, Y$ .

### Theorem 5.3.2

Let  $X, Y, Z$  be graphs. If  $f : Z \rightarrow X$  and  $g : Z \rightarrow Y$  are homomorphisms, there is a unique homomorphism  $\phi : Z \rightarrow X \times Y$  such that

$$\begin{aligned} f &= p_X \circ \phi \\ g &= p_Y \circ \phi. \end{aligned}$$

**Proof**

Define  $\phi(z) := (f(z), g(z))$  for all  $z \in Z$ . If  $u \sim v \in Z$ , then  $f(u) \sim f(v)$  in  $X$  and  $g(u) \sim g(v)$  in  $Y$ . Hence

$$(f(u), g(u)) \sim (f(v), g(v))$$

in  $X \times Y$  by construction. By definition,  $\phi$  is a homomorphism such that  $f = p_X \circ \phi, g = p_Y \circ \phi$ .

Suppose now that  $\phi(z) = (\phi_X(z), \phi_Y(z))$  is a function such that  $p_X \circ \phi(z) = \phi_X(z) = f(z)$  and  $p_Y \circ \phi(z) = \phi_Y(z) = g(z)$  for all  $z \in Z$ . But then  $\phi = \phi_X \times \phi_Y$  is determined by  $f, g$ .

We write

$$\text{Hom}(X, Y)$$

to denote the set of homomorphisms from  $X$  to  $Y$ .

**Corollary 5.3.2.1**

$$|\text{Hom}(Z, X \times Y)| = |\text{Hom}(Z, X)| \cdot |\text{Hom}(Z, Y)|.$$

**Proof**

By the previous lemma, there is an injective map

$$\text{Hom}(Z, X) \times \text{Hom}(Z, Y) \rightarrow \text{Hom}(Z, X \times Y).$$

Thus  $|\text{Hom}(Z, X \times Y)| \leq |\text{Hom}(Z, X)| \cdot |\text{Hom}(Z, Y)|$ .

For each  $\phi \in \text{Hom}(Z, X \times Y)$   $p_X \circ \phi \in \text{Hom}(Z, X)$  and  $p_Y \circ \phi \in \text{Hom}(Z, Y)$ . Thus we have an injective map

$$\text{Hom}(Z, X \times Y) \rightarrow \text{Hom}(Z, X) \times \text{Hom}(Z, Y).$$

Thus  $|\text{Hom}(Z, X \times Y)| \geq |\text{Hom}(Z, X)| \cdot |\text{Hom}(Z, Y)|$ .

**Definition 5.3.2 (Lattice)**

A partially ordered set  $\Lambda$  is a lattice, if for each  $x, y \in \Lambda$ , there is a least upper bound  $x \vee y$ , as well as a greatest lower bound  $x \wedge y$ .

**Lemma 5.3.3**

The set of isomorphism classes of cores with partial order  $\rightarrow$  is a lattice.

**Proof**

Least Upper Bound: Let  $X, Y$  be cores. We claim that  $X \vee Y = (X + Y)^\bullet$ .

It is clear that  $X \rightarrow X + Y \rightarrow (X + Y)^\bullet$  as well as  $Y \rightarrow X + Y \rightarrow (X + Y)^\bullet$ . Thus  $(X + Y)^\bullet$  is an upper bound of  $X \times Y$ .

Suppose now that  $Z$  is a core such that  $X \rightarrow Z, Y \rightarrow Z$ . Then  $X + Y \rightarrow Z$ .

Naturally,  $(X + Y)^\bullet \rightarrow Z$  through the identity map and hence  $X \vee Y = (X + Y)^\bullet$ .

Greatest Lower Bound: We show that  $X \wedge Y = (X \times Y)^\bullet$ .

Indeed,  $(X \times Y)^\bullet \rightarrow X \times Y \rightarrow X$  and  $(X \times Y)^\bullet \rightarrow X \times Y \rightarrow Y$ .

Suppose  $Z$  is a core such that  $Z \rightarrow X, Z \rightarrow Y$ . Then  $Z \rightarrow (X \times Y) \rightarrow (X \times Y)^\bullet$ . Hence  $X \wedge Y = (X \times Y)^\bullet$  as desired.

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# Chapter 6

## Matrix Theory

### 6.1 Spectral Graph Theory

#### 6.1.1 The Adjacency Matrix

**Definition 6.1.1 (Adjacency Matrix)**

Let  $X$  be an undirected graph. Its adjacency matrix  $A = A(X) \in \{0, 1\}^{V \times V}$  is the square matrix such that

$$A_{u,v} = 1\{u \sim v\}.$$

Note that  $A(X)$  is symmetric.

**Definition 6.1.2 (Graph Characteristic Polynomial)**

The characteristic polynomial of  $X$  is

$$\phi(A, x) := \det(xI - A(X)).$$

Note that characteristic polynomials do not uniquely determine  $X$ .

**Definition 6.1.3 (Graph Spectrum)**

The spectrum of  $X$  is the list of eigenvalues of  $A(X)$  together with their algebraic multiplicities.

If  $X \cong Y$ , then  $X, Y$  have the same spectrum.

**Lemma 6.1.1**

Let  $A = A(X)$ . Then  $(A^r)_{u,v}$  is the number of walks from  $u$  to  $v$  in  $X$  of length  $r$ .

**Corollary 6.1.1.1**

Suppose  $X$  has  $m$  edges and  $t$  triangles. Let  $A = A(X)$ . Then

$$\begin{aligned}\operatorname{tr} A &= 0 \\ \operatorname{tr} A^2 &= 2m \\ \operatorname{tr} A^3 &= 6t.\end{aligned}$$

Let  $n := |V(X)|$ . Remark that the number of closed  $v$ -walks is

$$\operatorname{tr} A^r = \sum_{i=1}^n \lambda_i^r.$$

**6.1.2 Incidence Matrix****Definition 6.1.4 (Incidence Matrix)**

The incidence matrix  $B = B(X) \in \{0, 1\}^{V \times E}$  of  $X$  defined by

$$(B)_{ue} = 1\{u \in e\}.$$

**Theorem 6.1.2**

Let  $X$  be a graph on  $n$  vertices and  $c$  bipartite components. Let  $B = B(X)$ . Then  $\operatorname{rank} B = n - c$ .

**Proof**

Remark that

$$\operatorname{rank} B = \operatorname{rank} B^T = n - \operatorname{null} B^T.$$

Hence it suffices to show that  $\operatorname{null} B^T = c$ .

Let  $z \in \mathbb{R}^n$  such that  $B^T z = 0$ . By the definition of  $B$ ,

$$z_u + z_v = 0$$

for all  $u, v$  such that  $u \sim v$ . In particular,

$$z_u = (-1)^r z_v$$

if  $u, v$  are joined by a path of length  $r$ .

It follows that  $z_u = 0$  if  $u$  is in a non-bipartite component and  $z$  takes inverse values on vertices from opposite classes in a bipartite component.

Thus  $\text{null}(B^T) = c$  as desired.

**Lemma 6.1.3**

Let  $B = B(X)$  and  $L$  the line graph of  $X$ . Then

$$B^T B = 2I + A(L).$$

**Definition 6.1.5 (Degree Matrix)**

The degree matrix of  $X$ ,  $D = D(X) \in \mathbb{Z}_+^{V \times V}$ , is the diagonal matrix where

$$(D)_{uu} = \text{deg}(u).$$

for each  $u \in V(X)$ .

**Lemma 6.1.4**

Let  $B = B(X)$ ,  $D = D(X)$ , and  $A = A(X)$ . Then

$$BB^T = D(X) + A(X).$$

## 6.2 Symmetric Matrices

### Proposition 6.2.1

Let  $A$  be a real symmetric  $n \times n$  matrix.

- (a) If  $u, v$  are eigenvectors of different eigenvalues, then  $u \perp v$ .
- (b) All eigenvalues are real.
- (c) If  $U$  is a subspace of  $\mathbb{R}^n$ , then  $U$  is  $A$ -invariant implies  $U^\perp$  is  $A$ -invariant.
- (d)  $U$  is a non-zero  $A$ -invariant subspace of  $\mathbb{R}^n$  implies that  $U$  contains a real eigenvector of  $A$ .
- (e)  $\mathbb{R}^n$  has an orthonormal basis consisting of eigenvectors of  $A$ .
- (f)  $A = PDP^T$ , where  $P$  is orthogonal, and columns of  $P$  are orthonormal eigenvectors of  $A$ .
- (g)  $A = \sum_{i=1}^n \lambda_i v_i v_i^T$  where  $\lambda_i$  is the eigenvalue corresponding to the orthonormal eigenvector  $v_i$ .

## 6.3 Eigenvectors of the Adjacency Matrix

We wish to find the eigenvalues of  $A = A(X)$ . Suppose there is some  $f : V(X) \rightarrow \mathbb{R}$  such that  $Af = \lambda f$ . By the definition of  $A$ ,

$$\begin{aligned}(Af)(u) &= \sum_v A_{uv} f(v) \\ &= \sum_{v \sim u} f(v).\end{aligned}$$

Hence if we can find (a vector/function)  $f$  such that

$$\sum_{v \sim u} f(v) = \lambda f(u)$$

for all  $u \in V(X)$ , Then  $\lambda$  is by definition an eigenvalue of  $A(X)$ .

### Example 6.3.1 (Regular Graphs)

For a  $r$ -regular graph  $X$ , by setting  $f = 1_n$ ,

$$\sum_{v \sim u} f(v) = r$$

Hence  $1_n$  is an eigenvector of  $X$ .

Note that the converse is also true.  $1_n$  is an eigenvector if and only if  $X$  is regular.

**Example 6.3.2 (Cycles)**

Let  $\tau$  be an  $n$ -th root of unity. Then let  $f(u) := \tau^u$ .

$$\sum_{v \sim u} f(v) = (\tau^{-1} + \tau)\tau^u$$

so  $\tau^{-1} + \tau$  is a real eigenvalue of  $X$ .

**Lemma 6.3.3**

Let  $X$  be  $k$ -regular with  $n$  vertices and eigenvalues

$$k, \theta_2, \dots, \theta_n.$$

Then  $X, \bar{X}$  have the same eigenvectors and the eigenvalues of  $\bar{X}$  are

$$n - k - 1, -\theta_2 - 1, \dots, -\theta_n - 1.$$

**Proof**

Let  $J$  be the all 1' matrix. By observation,

$$A(\bar{X}) = J - I - A(X).$$

Recall that  $1_n$  is the eigenvector of  $A(X)$  corresponding to eigenvalue  $k$ . Hence  $1_n$  is the eigenvector of  $A(\bar{X})$  corresponding to the eigenvalue  $n - 1 - k$ .

Let  $\{1_n, v_2, \dots, v_n\}$  be the orthogonal eigenvectors of  $A$ . For each  $2 \leq j \leq n$ ,

$$\begin{aligned} A(X)v_j &= \theta_j v_j \\ 1^T v_j &= 0. \end{aligned}$$

It follows that

$$\begin{aligned} A(\bar{X})v_j &= (J - I - A(X))v_j \\ &= -v_j - \theta_j v_j \\ &= (-1 - \theta_j)v_j. \end{aligned}$$

## 6.4 Positive Semidefinite Matrices

### Definition 6.4.1 (Positive Semidefinite Matrix)

A real symmetric matrix  $A$  is positive semidefinite (PSD) if for all  $u \in \mathbb{R}^n$ ,

$$u^T A u \geq 0.$$

If in addition,  $u^T A u = 0 \iff u = 0$ , then  $A$  is *positive definite*.

### Proposition 6.4.1

Let  $A \in \mathbb{R}^{n \times n}$ . The following are equivalent:

- (a)  $A$  is PSD
- (b) All eigenvalues of  $A$  are non-negative
- (c)  $A = B^T B$  for some  $B$

### Lemma 6.4.2

If  $L$  is a line graph, then  $\lambda_{\min}(L) \geq -2$ .

### Proof

Suppose  $L$  is the line graph of  $X$ . Let  $B = B(X)$ . We know that  $B^T B = A(L) + 2I$ . But  $B^T B$  is PSD. Hence  $A(L) + 2I$  has minimum eigenvalue 0 and  $\lambda_{\min}(L) \geq -2$ .

### Lemma 6.4.3

Let  $Y$  be an induced subgraph of  $X$ . Then

$$\lambda_{\min}(X) \leq \lambda_{\min}(Y) \leq \lambda_{\max}(Y) \leq \lambda_{\max}(X).$$

### Proof

Let  $A = A(X)$  and  $\tilde{A} = A(Y)$ . Put  $\lambda = \lambda_{\max}(X)$ . Thus  $\lambda I - A$  is PSD.

For any  $\tilde{f} : V(Y) \rightarrow \mathbb{R}$ , extend it to a function  $f : V(X) \rightarrow \mathbb{R}$  such that  $f(u) = 0$  for each  $u \in V(X) \setminus V(Y)$ . Then

$$\begin{aligned} 0 &\leq f^T (\lambda I - A) f \\ &\leq \tilde{f}^T (\lambda I - \tilde{A}) \tilde{f}. \end{aligned}$$

Thus  $\lambda I - \tilde{A}$  is PSD and  $\lambda_{\max}(\tilde{A}) \leq \lambda$ .

Similarly, working on the PSD matrix  $A(X) - \lambda_{\min}(X) \cdot I$ , we can show  $\lambda_{\min}(\tilde{A}) \geq \lambda_{\min}(A)$ .

Remark that this lemma follows from a more general interlacing theorem.

**Definition 6.4.2 (Laplacian Matrix)**

The Laplacian matrix of the graph  $X$ , is

$$L(X) := D(X) - A(X).$$

**Proposition 6.4.4**

$L := L(X)$  is positive semidefinite.

**Proof**

Let  $n := |V(X)|$ . For any  $v \in \mathbb{R}^n$ ,

$$\begin{aligned} x^T Lx &= \sum_{u,v} x_u L_{uv} x_v \\ &= \sum_u x_u^2 \deg(u) - \sum_u x_u \sum_{v \sim u} x_v \\ &= \sum_{uv \in E} (x_u^2 + x_v^2) - \sum_{uv \in E} 2x_u x_v \\ &= \sum_{uv \in E} (x_u - x_v)^2 \\ &\geq 0. \end{aligned}$$

Remark that  $x^T Lx$  measures the “smoothness” of  $x$  on  $X$ .

Since  $L(X)$  is PSD, its smallest eigenvalue is at least 0. Observe that

$$(L(X)1^n)_u = \deg(u) - \deg_u = 0.$$

Hence  $1_n$  is an eigenvector of  $L(X)$  with eigenvalue 0 and the minimum eigenvalue of  $L(X)$  is 0.

**Proposition 6.4.5**

Let  $L = L(X)$  and

$$0 = \mu_1 \leq \mu_2 \leq \dots \leq \mu_n$$

be the eigenvalues of  $L$ . Then  $\mu_2 > 0$  if and only if  $X$  is connected.

**Proof**

( $\implies$ ) Suppose  $X$  is the disjoint union of 2 graphs  $X_1, X_2$  with  $n_1, n_2$  vertices. Then

$$L = \begin{bmatrix} L(X_1) & 0 \\ 0 & L(X_2) \end{bmatrix}.$$

Both  $\begin{bmatrix} 0_{n_1} \\ 1_{n_2} \end{bmatrix}$  and  $\begin{bmatrix} 1_{n_1} \\ 0_{n_2} \end{bmatrix}$  are eigenvectors with eigenvalue 0 and  $\mu_2 = 0$ .

( $\impliedby$ ) Let  $f$  be an eigenvector with eigenvalue 0. But then

$$\begin{aligned} 0 &= f^T(Lf) \\ &= \sum_{uv \in E(X)} (f(u) - f(v))^2 \\ &\iff \forall uv \in E(X), f(u) = f(v). \end{aligned}$$

Since  $X$  is connected,  $f$  is constant on  $V(X)$ . The eigenspace corresponding to eigenvalue 0 has dimension 1. It follows that  $\mu_2 > 0$ .

Suppose  $X$  is  $k$ -regular. Let  $A = A(X)$ . Then

$$L = L(X) = kI - A(X).$$

If  $k = \lambda_1 \geq \dots \geq \lambda_n$  are the eigenvalues of  $A$  and  $0 = \mu_1 \leq \dots \leq \mu_n$  are the eigenvalues of  $L$ , then  $\mu_i = k - \lambda_i$  for each  $1 \leq i \leq n$ .

## 6.5 Courant-Fisher Theorem

### Definition 6.5.1 (Rayleigh Quotient)

The Rayleigh quotient of  $X \in \mathbb{R}^n$  with respect to  $A \in \mathbb{R}^{n \times n}$  is

$$\frac{x^T Ax}{x^T x}.$$

Observe that if  $x$  is an eigenvector of  $A$  with respect to the eigenvalue  $\lambda$ ,

$$\frac{x^T Ax}{x^T x} = \lambda.$$

The Courant-Fisher theorem gives a relation between the extreme values of the Rayleigh quotients and the eigenvalues of  $A$ .



**Theorem 6.5.1 (Courant-Fisher)**

Let  $A \in \mathbb{R}^{n \times n}$  be symmetric with eigenvalues

$$\lambda_1 \geq \lambda_2 \geq \cdots \geq \lambda_n.$$

Then for all  $k \in [n]$ ,

$$\lambda_k = \max_{S \subseteq \mathbb{R}^n: \dim S = k} \min_{x \in S: x \neq 0} \frac{x^T A x}{x^T x} = \min_{S \subseteq \mathbb{R}^n: \dim S = n - k + 1} \max_{x \in S: x \neq 0} \frac{x^T A x}{x^T x}.$$

**Proof**

Let  $f_1, \dots, f_n$  be an orthonormal eigenvector basis of  $A$  where  $f_i$  corresponds to  $\lambda_i$ . Consider  $S := \text{span}\{f_1, \dots, f_k\}$ .

We first show that

$$\max_{S \subseteq \mathbb{R}^n: \dim S = k} \min_{x \in S: x \neq 0} \frac{x^T A x}{x^T x} \geq \lambda_k.$$

For any  $x = \sum_{i=1}^k c_i f_i \in S$ ,

$$\begin{aligned} x^T A x &= \left( \sum_{i=1}^k c_i f_i \right)^T A \left( \sum_{i=1}^k c_i f_i \right) \\ &= \left( \sum_{i=1}^k c_i f_i \right)^T \left( \sum_{i=1}^k c_i \lambda_i f_i \right) \\ &= \sum_{i=1}^k \lambda_i c_i^2. \end{aligned}$$

Moreover,

$$x^T x = \sum_{i=1}^k c_i^2.$$

Thus the Rayleigh quotient of  $x$  is

$$\begin{aligned} \frac{x^T A x}{x^T x} &= \frac{\sum_{i=1}^k \lambda_i c_i^2}{\sum_{i=1}^k c_i^2} \\ &\geq \lambda_k \frac{\sum_{i=1}^k c_i^2}{\sum_{i=1}^k c_i^2} \\ &= \lambda_k. \end{aligned}$$

We then show that

$$\max_{S \subseteq \mathbb{R}^n: \dim S = k} \min_{x \in S: x \neq 0} \frac{x^T A x}{x^T x} \leq \lambda_k.$$

Let  $S$  be a subspace of dimension  $k$  and  $T := \text{span}\{f_k, \dots, f_n\}$  be a subspace of dimension  $n - k + 1$ .

But then by dimension,  $\dim(S \cap T) \geq 1$ .

$$\begin{aligned} \min_{x \in S: x \neq 0} \frac{x^T A x}{x^T x} &\leq \min_{x \in S \cap T: x \neq 0} \frac{x^T A x}{x^T x} \\ &\leq \max_{x \in S \cap T: x \neq 0} \frac{x^T A x}{x^T x} \\ &\leq \max_{x \in T: x \neq 0} \frac{x^T A x}{x^T x} \\ &\leq \lambda_k. \end{aligned}$$

The last inequality follows from an identical calculation to our work above.

The other equality is identical and we omit the proof.

### Theorem 6.5.2

Let  $A \in \mathbb{R}^{n \times n}$  be symmetric and  $x \neq 0$  a vector that maximizes or minimizes the Rayleigh quotient. Then  $x$  is an eigenvector of  $\lambda_1$  or  $\lambda_n$  respectively.

## 6.6 Maximum Graph Eigenvalue

### Theorem 6.6.1

Let  $\bar{d}, \Delta$ , be the average and maximum degrees of  $x$ . Then

$$\bar{d} \leq \lambda_1 \leq \Delta.$$

**Proof**

$\bar{d} \leq \lambda_1$ : By the Courant-Fisher theorem,

$$\begin{aligned}\lambda_1 &= \max_{x \in \mathbb{R}^n: x \neq 0} \frac{x^T A x}{x^T x} \\ &\geq \frac{1^T A 1}{1^T 1} \\ &= \frac{\sum_u \sum_{v \sim u} 1}{n} \\ &= \frac{2|E(X)|}{n} \\ &= \bar{d}.\end{aligned}$$

$\lambda_1 \leq \Delta$ : Let  $f$  be the unit eigenvector corresponding to  $\lambda_1$ . Let  $u := \operatorname{argmax}_{v \in V} f(v)$ . Without loss of generality, by taking  $-f$  if necessary, we may assume that  $f(u) > 0$ .

Then

$$\begin{aligned}\lambda_1 &= \frac{(Af)(u)}{f(u)} \\ &= \frac{\sum_{v \sim u} f(v)}{f(u)} \\ &\leq \frac{\deg(u) \cdot f(u)}{f(u)} \\ &\leq \Delta\end{aligned}$$

as desired.

**Lemma 6.6.2**

If  $X$  is connected with  $\lambda_1 = \Delta$ , then  $X$  is  $\Delta$ -regular.

**Proof**

Let  $f$  be the unit eigenvector corresponding to  $\lambda_1$ . Let  $u := \operatorname{argmax}_{v \in V} f(v)$ . Without loss of generality, by taking  $-f$  if necessary, we may assume that  $f(u) > 0$ .

Then

$$\begin{aligned}\lambda_1 &= \frac{(Af)(u)}{f(u)} \\ &= \frac{\sum_{v \sim u} f(v)}{f(u)} \\ &\leq \frac{\deg(u) \cdot f(u)}{f(u)} \\ &\leq \Delta \\ &= \lambda_1.\end{aligned}$$

Now this holds with equality if and only if  $f(v) = f(u)$  for all  $v \sim u$  and  $\deg(u) = \Delta$ .

By connectivity, we can apply the same argument in breadth-first search fashion to realize that  $f$  is constant on  $V(X)$ . It must then be that each  $v \in V$  satisfies  $\deg(v) = \Delta$  as we argued for  $u$  above.

**Theorem 6.6.3 (Wilf '67)**

The chromatic number of a graph  $X$  satisfies

$$\chi(X) \leq \lfloor \lambda_1 \rfloor + 1.$$

**Proof**

We argue by induction on  $n = |V(X)|$ . The base case of  $n = 1$  is trivial as  $\chi(X) = 1, \lambda_1 = 0$ . So the theorem is true.

Suppose  $n \geq 1$  and the theorem holds for all graphs with  $n$  vertices. Let  $X$  be a graph on  $n + 1$  vertices and  $\lambda_1 := \lambda_1(X)$ . By a previous lemma, the average degree of  $X$  is at most  $\lambda_1$ . Hence there is at least one vertex  $u$  for which  $\deg(u) \leq \lfloor \lambda_1 \rfloor$ . Define  $S := V(X - u)$ .

Let  $\lambda = \lambda_1(X[S])$ . By our work earlier,  $\lambda \leq \lambda_1(X) = \lambda_1$  since  $X[S]$  is an induced subgraph of  $X$ . By the induction hypothesis,

$$\chi(X[S]) \leq \lfloor \lambda \rfloor + 1 \leq \lfloor \lambda_1 \rfloor + 1.$$

Consider any  $(\lfloor \lambda_1 \rfloor + 1)$ -colouring for  $X[S]$ . We can extend this to a colouring of  $X$  since  $u$  has at most  $\lfloor \lambda_1 \rfloor$  neighbours. By induction, the theorem holds.

## 6.7 Eigenvalue Interlacing

### Definition 6.7.1 (Principal Submatrix)

A submatrix  $B$  of  $A \in \mathbb{R}^{n \times n}$  is principal if  $B$  is obtained by deleting the same set of rows and columns of  $A$ .

### Theorem 6.7.1 (Cauchy)

Let  $A \in \mathbb{R}^{n \times n}$  be symmetric and  $B$  an  $(n-1) \times (n-1)$  principal submatrix of  $A$ . Let  $\alpha_1 \geq \dots \geq \alpha_n$  and  $\beta_1 \geq \dots \geq \beta_{n-1}$  be the spectra of  $A, B$  respectively.

Then

$$\alpha_1 \geq \beta_1 \geq \alpha_2 \geq \beta_2 \geq \dots \geq \alpha_{n-1} \geq \beta_{n-1} \geq \alpha_n.$$

### Proof

Without loss of generality, assume that  $B$  is obtained from  $A$  by deleting the first row and column. By the Courant-Fisher theorem,

$$\begin{aligned} \alpha_k &= \max_{S \subseteq \mathbb{R}^n: \dim S = k} \min_{x \in S: x \neq 0} \frac{x^T A x}{x^T x} \\ \beta_k &= \max_{S \subseteq \mathbb{R}^{n-1}: \dim S = k} \min_{x \in S: x \neq 0} \frac{x^T B x}{x^T x} \\ &= \max_{S \subseteq \mathbb{R}^{n-1}: \dim S = k} \min_{x \in S: x \neq 0} \frac{\begin{bmatrix} 0 \\ x \end{bmatrix}^T A \begin{bmatrix} 0 \\ x \end{bmatrix}}{\begin{bmatrix} 0 \\ x \end{bmatrix}^T \begin{bmatrix} 0 \\ x \end{bmatrix}} \\ &= \max_{S \subseteq \mathbb{R}^{n-1}: \dim S = k, S \perp e_1} \min_{x \in S: x \neq 0} \frac{x^T A x}{x^T x} \\ &\leq \alpha_k. \end{aligned}$$

With an identical argument applied to  $-A, -B$ , we get that  $-\beta_k \leq -\alpha_{k+1}$  and hence

$$\beta_k \geq \alpha_{k+1}.$$

## 6.7.1 Perron-Frobenius Theorem

### Theorem 6.7.2 (Perron-Frobenius; Undirected Graphs)

Let  $X$  be connected and  $A := A(X)$ . Let  $\lambda_1 \geq \dots \geq \lambda_n$  be the spectrum of  $A$ . Then

- (a)  $\lambda_1$  has a strictly positive eigenvector
- (b)  $\lambda_1 > \lambda_2$
- (c)  $\lambda_1 \geq -\lambda_n$  with equality if and only if  $X$  is bipartite.

#### Proof (a)

We first show that  $\lambda_1$  has a non-negative eigenvector. Let  $f$  be a unit eigenvector of  $\lambda_1$  and consider

$$f^+ := |f|.$$

We have

$$\begin{aligned}\lambda_1 &= f^T A f \\ &= \sum_{uv \in E(X)} 2f(u)f(v) \\ &\leq \sum_{uv \in E(X)} 2f^+(u)f^+(v) \\ &= (f^+)^T A f^+.\end{aligned}$$

But we must have equality since  $f$  is a maximizer of the Rayleigh quotient. Hence  $f^+$  is also a maximizer. By a previous theorem, this implies that  $f^+$  is also an eigenvector of  $\lambda_1$ .

Now suppose there is some  $u$  for which  $f^+(u) = 0$ . Since  $f \not\equiv 0$  and  $X$  is connected, there is some  $vw \in E(X)$  for which  $f^+(v) = 0 < f^+(w)$ . But then

$$\sum_{w \sim v} f^+(w) > 0 = \lambda_1 f^+(v)$$

which is a contradiction.

Hence  $f^+$  is a strictly positive eigenvector of  $\lambda_1$ .

**Proof (b)**

Let  $g$  be a unit eigenvector of  $\lambda_2$  which is orthogonal to  $f^+$ . Hence

$$\begin{aligned}
\lambda_1 &= g^T A g \\
&= \sum_{vw \in E(X)} g(v)g(w) \\
&\leq \sum_{vw \in E(X)} |g(v)| |g(w)| \\
&= |g|^T A |g| \\
&\leq \max_{g \in \mathbb{R}^n: \|g\|=1} g^T A g \\
&= \lambda_1.
\end{aligned}$$

If we have  $\lambda_1 = \lambda_1$  then  $|g|$  is strictly positive by the same argument as in the proof of (a). Moreover, for each  $uv \in E(X)$ ,

$$g(u)g(v) = |g(u)||g(v)|.$$

Now,  $f^+$  is strictly positive,  $g^T f^+ = 0$ , and  $g$  has non-zero entries. Hence there must be two entries of  $g$  at which the signs differ. But since  $X$  is connected, there is an edge  $uv \in E(X)$  at which  $g(u) < 0 < g(v)$ . But then

$$g(u)g(v) \neq |g(u)||g(v)|$$

which contradicts our conclusion above.

It follows that  $\lambda_2 < \lambda_1$ .

**Proof (c)**

Let  $g$  be a unit eigenvector of  $\lambda_n$ . Then

$$\begin{aligned}
|\lambda_n| &= |g^T A g| \\
&= \left| \sum_{uv \in E(X)} 2g(u)g(v) \right| \\
&\leq \sum_{uv \in E(X)} 2|g(u)| \cdot |g(v)| \\
&= |g|^T A |g| \\
&\leq \max_{g \in \mathbb{R}^n: \|g\|=1} g^T A g \\
&= \lambda_1.
\end{aligned}$$

$X$  bipartite implies  $\lambda_n = -\lambda_1$ . Let  $S, T$  be a bipartition of  $X$ . Suppose  $X$  has eigenvalue  $\lambda$  with eigenvector  $f$ . Define

$$\bar{f} := \begin{cases} f(u), & u \in S \\ -f(u), & u \in T \end{cases}$$

Then for  $u \in S$ ,

$$\begin{aligned} (A\bar{f})(u) &= \sum_{v \sim u} -f(v) \\ &= -\lambda f(u) \\ &= -\bar{f}(u). \end{aligned}$$

For  $u \in T$ ,

$$\begin{aligned} (A\bar{f})(u) &= \sum_{v \sim u} f(v) \\ &= -\lambda \bar{f}(u). \end{aligned}$$

Thus  $-\lambda$  is an eigenvalue of  $X$  as well. It follows that the spectrum of  $X$  is symmetric around 0. In particular,  $\lambda_1 = \lambda_n$ .

$\lambda_1 = \lambda_n$  implies  $X$  is bipartite. If  $|\lambda_n| = \lambda_1$ , then  $|g|$  is a strictly positive eigenvector of  $\lambda_1$ . Moreover, for equality to hold in the initial inequality above, every edge  $uv \in E(X)$  satisfies

$$g(u)g(v) = -|g(u)||g(v)| < 0.$$

Let  $S$  be the vertices  $v$  for which  $g(v) > 0$  and  $T$  the vertices  $v$  for which  $g(v) < 0$ . Both are non-empty by our work above.

It follows that all edges in  $X$  joins  $S$  and  $T$  so it forms a bipartition of  $X$  as desired.



# Chapter 7

## Strongly Regular Graphs

### 7.1 Parameters

**Definition 7.1.1 (Strongly Regular)**

$X$  is strongly regular with parameters  $(n, k, a, c)$  if

- (i)  $|V(X)| = n$
- (ii)  $X$  is  $k$ -regular
- (iii) Every pair of adjacent vertices has  $a$  common neighbours
- (iv) Every pair of non-adjacent vertices has  $c$  common neighbours

**Example 7.1.1**

$C_5$  is a  $(5, 2, 0, 1)$ -strongly regular graph.

**Proposition 7.1.2**

Suppose  $X$  is  $(n, k, a, c)$  strongly regular. Then  $\bar{X}$  is  $(n, \bar{k}, \bar{a}, \bar{c})$  strongly regular for

$$\begin{aligned}\bar{k} &= n - 1 - k \\ \bar{a} &= n - 2 - 2k + c \\ \bar{c} &= n - 2k + a\end{aligned}$$

**Definition 7.1.2 (Primitive)**

A strongly regular graph  $X$  is primitive if both  $X, \bar{X}$  are connected.

If  $X$  is not primitive, then it is *imprimitive*.

**Lemma 7.1.3**

Let  $X$  be  $(n, k, a, c)$  strongly regular. Then the following are equivalent:

- (a)  $X$  is disconnected
- (b)  $c = 0$
- (c)  $a = k - 1$
- (d)  $X \cong mK_{k-1}$ . That is, it is isomorphic to  $m$  copies of  $K_{k-1}$  for some  $m$

**Proof**

(a)  $\implies$  (b) Suppose  $u, v$  are in different components. Then  $N(u) \cap N(v) = \emptyset$  implies  $c = 0$ .

(b)  $\implies$  (c) Fix any  $u \in V(X)$ . We claim that all neighbours of  $u$  must be pairwise adjacent. Let  $v, w \in N(u)$ . Then  $u \in N(v) \cap N(w)$ . Since  $c = 0$ , it must be that  $v \sim w$ .

Thus  $a = k - 1$ .

(c)  $\implies$  (d), (d)  $\implies$  (a) Obvious.

**Example 7.1.4**

The line graph of  $K_n$  is

$$\left( \binom{n}{2}, 2n - 4, n - 2, 4 \right)$$

strongly regular.

The line graph of  $K_{n,n}$  is

$$(n^2, 2n - 2, n - 2, 2)$$

strongly regular.

### 7.1.1 Parameter Relation

**Proposition 7.1.5**

A  $(n, k, a, c)$  strongly regular graph  $X$  satisfies

$$c(n - k - 1) = k(k - a - 1).$$

**Proof**

Fix  $u \in V(X)$ . Let  $U$  denote the neighbours of  $u$  and  $V$  the non-neighbours of  $u$ .

Clearly,  $|U| = k$ ,  $|V| = n - k - 1$ .

Now, for each  $v \in V$ ,  $v \not\sim u$  so they must share exactly  $c$  neighbours. But all neighbours of  $u$  are in  $U$  and

$$E(U, V) = c(n - k - 1).$$

Similarly, for all  $x \in U$ ,  $x \sim u$  hence they share  $a$  common neighbours. Since all neighbours of  $u$  are in  $U$ ,  $x$  shares  $a$  neighbours with  $u$  in  $U$ . Then  $x$  has  $k - 1 - a$  other neighbours in  $V$ . Hence

$$E(U, V) = k(k - 1 - a).$$

## 7.2 Eigenvalues

**Lemma 7.2.1**

Suppose  $X$  is  $(n, k, a, c)$  strongly regular and  $X$  is not complete. Put  $A := A(X)$ . Then  $A$  has 3 distinct eigenvalues  $k, \theta, \tau$ , where

$$\begin{aligned}\theta &= \frac{a - c + \sqrt{\Delta}}{2} \\ \tau &= \frac{a - c - \sqrt{\Delta}}{2} \\ \Delta &:= (a - c)^2 + 4(k - c).\end{aligned}$$

These are the roots of the quadratic polynomial

$$\lambda^2 - (a - c)\lambda - (k - c).$$

Moreover, the algebraic multiplicities of  $\theta, \tau$  are

$$\begin{aligned}m_\theta &= \frac{1}{2} \left( (n - 1) - \frac{2k + (n - 1)(a - c)}{\sqrt{\Delta}} \right) \\ m_\tau &= \frac{1}{2} \left( (n - 1) + \frac{2k + (n - 1)(a - c)}{\sqrt{\Delta}} \right).\end{aligned}$$

**Proof**

Consider the entries of  $A^2$ . There are  $k$  closed 2-walks at each vertex. If  $u \sim v$ , there are

precisely  $a$  2-walks from  $u$  to  $v$ . If  $u \not\sim v$ , there are exactly  $c$  2-walks from  $u$  to  $v$ .

It follows that

$$\begin{aligned} A^2 &= kI + aA + c(J - I - A) \\ A^2 - (a - c)A - (k - c)I &= cJ. \end{aligned}$$

We know that  $k$  is an eigenvalue of  $A$  with the all-1s vector  $1_n$  being an eigenvector. Thus if  $\lambda \neq k$  is an eigenvalue with eigenvector  $v$ , then  $v \perp 1_n$ .

Moreover,

$$\begin{aligned} A^2v - (a - c)Av - (k - c)Iv &= cJv \\ \lambda^2v - (a - c)\lambda v - (k - c)v &= 0 \\ \lambda^2 - (a - c)\lambda - (k - c) &= 0. \end{aligned}$$

Clearly  $\lambda \in \theta, \tau$ .

It remains to verify  $m_\theta, m_\tau$  as well as  $m_\theta, m_\tau > 0$  with  $\theta, \tau \neq k$  and  $\theta \neq \tau$ .

$\theta \neq \tau$  We have  $\Delta := (a - c)^2 + 4(k - c) = 0$  if and only if  $a = k = c$ . But since  $a \leq k - 1$ , this is impossible.

By the definitions of  $\theta, \tau$ , they must differ.

$\theta, \tau \neq k$  This follows directly from the Perron-Frobenius theorem.

$m_\theta, m_\tau$  There are  $n$  eigenvalues, hence

$$m_\theta + m_\tau = n - 1.$$

Moreover,  $\text{tr } A = 0$  implies that

$$k + m_\theta\theta + m_\tau\tau = 0.$$

Solving these linear equations yield

$$\begin{aligned} m_\theta &= -\frac{(n - 1)\tau + k}{\theta - \tau} \\ m_\tau &= \frac{(n - 1)\theta + k}{\theta - \tau} \end{aligned}$$

$m_\theta, m_\tau > 0$  We leave this as an exercise.

**Lemma 7.2.2**

A connected regular graph with exactly 3 distinct eigenvalues is strongly regular.

**Proof**

Suppose  $X$  is  $k$ -regular, connected, and has eigenvalues  $k, \theta, \tau$ . Put  $A := A(X)$  and

$$M := (k - \theta)^{-1}(k - \tau)^{-1}(A - \theta I)(A - \tau I).$$

Then  $M$  has eigenvalue 1 with multiplicity 1 by the Perron-Frobenius theorem with corresponding eigenvector  $1_n$ . Moreover, we can also show that  $M$  has eigenvalue 0 with multiplicity  $n - 1$ .

Now,  $\text{rank } M = 1$ . Moreover,  $M$  is symmetric and  $M1_n = 1_n$ . Thus we must have

$$M = \frac{1}{n}J.$$

By computation,

$$\begin{aligned} (A - \theta I)(A - \tau I) &= (k - \theta)(k - \tau)\frac{1}{n}J \\ A^2 - (\theta + \tau)A + \theta\tau I &= (k - \theta)(k - \tau)\frac{1}{n}J \\ A^2 &= (\theta + \tau)A - \theta\tau I + (k - \theta)(k - \tau)\frac{1}{n}J. \end{aligned}$$

It follows that the number of 2-walks from  $u$  to  $v$  in  $X$  depends only on whether  $u = v$ ,  $u \sim v$ , or  $u \neq v, u \not\sim v$ . Hence  $X$  is strongly regular.

**Theorem 7.2.3**

Suppose  $X$  is connected and regular and is NOT complete. Then  $X$  has 3 distinct eigenvalues if and only if  $X$  is strongly regular.

Recall that  $\Delta := (a - c)^2 + 4(k - c)$  and

$$\begin{aligned} m_\theta &= \frac{1}{2} \left( (n - 1) - \frac{2k + (n - 1)(a - c)}{\sqrt{\Delta}} \right) \\ m_\tau &= \frac{1}{2} \left( (n - 1) + \frac{2k + (n - 1)(a - c)}{\sqrt{\Delta}} \right). \end{aligned}$$

But  $m_\theta, m_\tau$  are integers. Hence this is a condition on the parameters  $n, k, a, c$ .

## 7.3 Payley Graphs

### Definition 7.3.1 (Payley Graphs)

Suppose  $q$  is a prime power and  $q \equiv 1 \pmod{4}$ . Then  $P(q)$  is the graph with

$$\begin{aligned}V(P(q)) &= \text{GF}(q) \\E(P(q)) &= \{uv : u - v \text{ is a nonzero square}\}.\end{aligned}$$

### Proposition 7.3.1

$P(q)$  is well-defined.

#### Proof

It suffices to show that  $u - v$  is a non-zero square if and only if  $v - u$  is a non-zero square.

It is not hard to see that exactly half of elements in  $\text{GF}(q) \setminus \{0\}$  are squares. In particular,  $P(q)$  is  $\frac{q-1}{2}$ -regular.

### Lemma 7.3.2

$P(q)$  is arc transitive.

#### Proof

Let  $\phi_{a,b} : V(P(q)) \rightarrow V(P(q))$  be given by

$$x \mapsto ax + b.$$

Then  $\phi_{a,b} \in \text{Aut}(P(q))$  if  $a$  is a non-zero square (exercise).

For any two arcs  $(x, y), (x', y')$ , let  $a = (x' - y')(x - y)^{-1}$  and  $b = x' - ax$ . It follows that

$$(x', y') = (x, y)^{\phi_{a,b}}$$

and  $P(q)$  is arc transitive.

### Theorem 7.3.3

$P(q)$  is self-complementary.

**Proof**

Define  $\sigma : x \mapsto ax$  where  $a \neq 0$  is NOT a square. This is an isomorphism to its complement.

**Theorem 7.3.4**

$P(q)$  is  $(q, \frac{q-1}{2}, \frac{q-5}{4}, \frac{q-1}{4})$  strongly regular.

**Proof**

Exercise.

## 7.4 More Strongly Regular Graphs

We have seen that  $C_n, P(q)$  are strongly regular.

**Definition 7.4.1 (Lattice Graph)**

Let  $L_n$  denote the graph where

$$\begin{aligned} V(L_n) &= [n] \times [n] \\ E(L_n) &= \{(a, b)(a', b') : a = a', b = b'\} \end{aligned}$$

Observe that  $L_n$  is  $(n^2, 2(n-1), n-2, 2)$  strongly regular.

**Definition 7.4.2 (Latin Square)**

An  $n \times n$  Latin square is an  $n \times n$  matrix such that each row and column is a permutation of  $[n]$ .

**Definition 7.4.3 (Latin Square Graph)**

Let  $A$  be an  $n \times n$  latin square. The Latin square graph  $X(A)$  associated with  $A_n$  is defined by

$$\begin{aligned} V(X(A)) &= [n] \times [n] \\ E(X(A)) &= \{(a, b)(a', b') : a = a' \vee b = b' \vee A_{ab} = A_{a'b'}\} \end{aligned}$$

The latin graph  $X(A)$  for an  $n \times n$  Latin square is  $(n^2, 3(n-1), n, 6)$  strongly regular (exercise).

## 7.5 The Second Eigenvalue

Let  $\lambda_2$  denote the second largest eigenvalue of  $A(X)$  and  $\mu_2$  denote the second smallest eigenvalue of  $L(X)$ . They are closely related. Indeed, if  $X$  is  $k$ -regular, then  $\mu_2 = k - \lambda_2$ .

While  $\lambda_1$  gives information about the degrees of  $X$ ,  $\lambda_2, \mu_2$  measure the expansion property of  $X$  (how well connected  $X$  is). We know that  $\mu_2 = 0$  if and only if  $X$  is disconnected. If  $X$  has a small  $\lambda_2$ , or large  $\mu_2$ , then  $X$  is a good expander.

### Definition 7.5.1 (Isoperimetric Ratio)

Let  $S \subseteq V(X)$ . The isoperimetric ratio of  $S$  is

$$\theta(S) := \frac{e(S, \bar{S})}{|S|}.$$

The isoperimetric ratio of  $X$  is

$$\theta(X) := \min_{S: |S| \leq \frac{V(X)}{2}} \theta(S).$$

### Definition 7.5.2 ( $\alpha$ -Expander)

$X$  is an  $\alpha$ -expander if  $\theta(X) \geq \alpha$ .

### Theorem 7.5.1

$$\theta(X) \geq \frac{\mu_2}{2}.$$

#### Proof

We argue that for every  $S \subseteq V(X), |S| \leq \frac{n}{2}$ , the isoperimetric ratio satisfies

$$\theta(S) \geq \mu_2 \left(1 - \frac{|S|}{n}\right).$$

The result then follows.

Recall that  $L = L(X)$  has  $\mu_1 = 0$  with eigenvector  $1_n$ . Hence  $\mu_2$  is the smallest eigenvalue of  $L - \frac{1}{n}J$ . Moreover, the Rayleigh quotient is minimized at an eigenvector  $v$  of  $\mu_2$  where  $v \perp 1_n$ .



That is,

$$\begin{aligned}\mu_2 &= \min_{x \perp 1_n} \frac{x^T (L - \frac{1}{n}J) x}{x^T x} \\ &= \min_{x \perp 1_n} \frac{x^T L x}{x^T x}.\end{aligned}$$

Let  $\alpha := \frac{|S|}{n}$  and  $x := \chi_S - \alpha 1_n$ . Then  $x \perp 1_n$  and

$$\|x\|_2^2 = \alpha n(1 - \alpha).$$

Moreover,

$$\begin{aligned}\mu_2 &\leq \frac{x^T L x}{x^T x} \\ &= \frac{1}{\alpha n(1 - \alpha)} \sum_{uv \in E} (x_u - x_v)^2 \\ &= \frac{1}{\alpha n(1 - \alpha)} e(S, \bar{S}) \\ &= \frac{\theta(S)}{1 - \alpha}.\end{aligned}$$